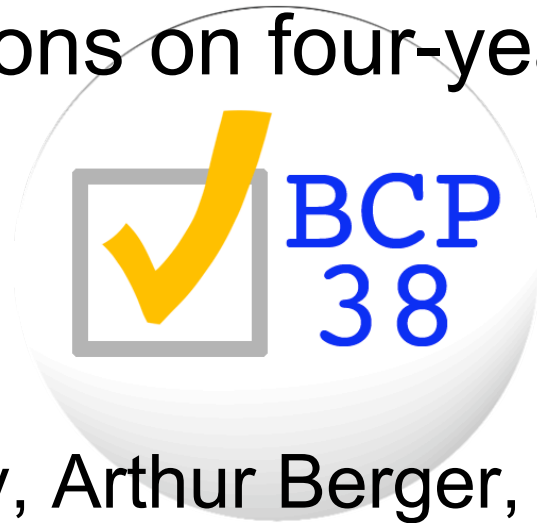


IP Spoofing Project

Observations on four-years of data



Rob Beverly, Arthur Berger, Young Hyun
{rbeverly,awberger}@csail.mit, youngh@caida

Spoofers Project

- Background
- Recent Relevance
- Project Description
- What's New: Methodology
- What's New: Data
- Parting Thoughts

Spoofed-Source IP Packets

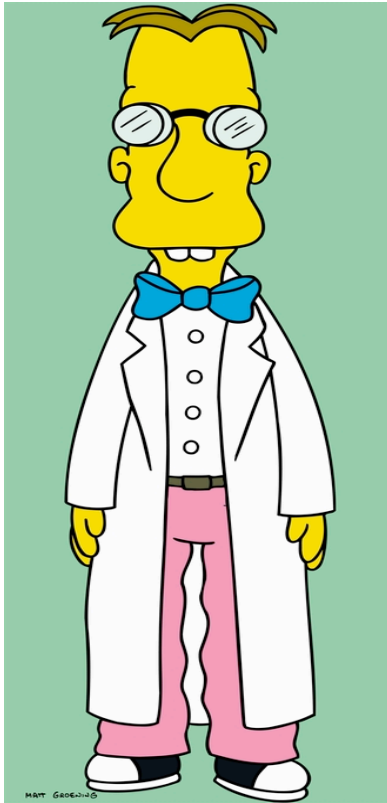
- Circumvent host network stack to forge or “spoof” *source address* of an IP packet
- Lack of source address accountability a basic Internet weakness:
 - Anonymity, indirection [VP01], amplification
- Security issue for more than two-decades [RTM85, Bellovin89]
- Still an attack vector?

0	4	8	16	19	31
Version	HLen	Tos	Length		
Ident		Flags	Offset		
TTL	Protocol	Checksum			
Source Address					
Destination Address					
Options (Variable)				Padding (Variable)	
Data					

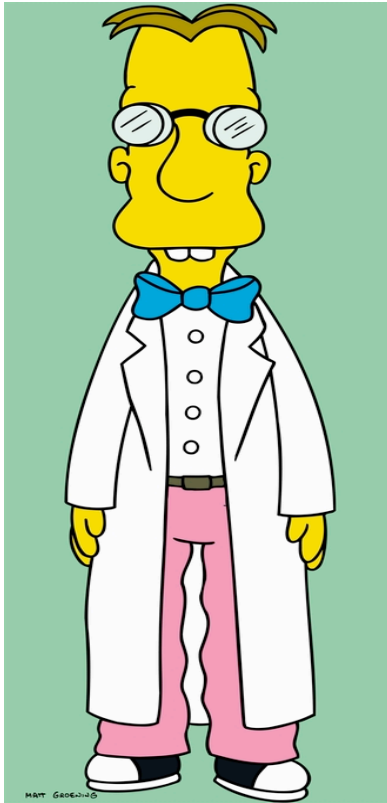
Circa 2004...

IP Source Spoofing doesn't matter!

- a) All providers **filter**
- b) All modern attacks use **botnets**
- c) Compromised hosts are behind **NATs**



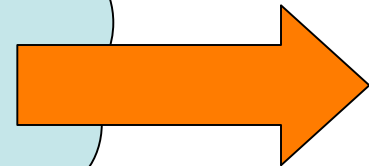
Circa 2004...



IP Source Spoofing doesn't matter!

- a) All providers **filter**
- b) All modern attacks use **botnets**
- c) Compromised hosts are behind **NATs**

!?!?!?



The Spoofer Project

- Strong opinions from many sides:
 - Academic
 - Operational
 - Regulatory
- ...but only anecdotal data

spoofer.csail.mit.edu

- Internet-wide active measurement effort:
 - Quantify the extent and nature of Internet source address filtering
- We learn and form inferences over:
 - Filtering policies/currently employed defenses
 - Filtering specificity, locations, providers, etc.
 - Distribution of filtering
- Began Feb. 2005



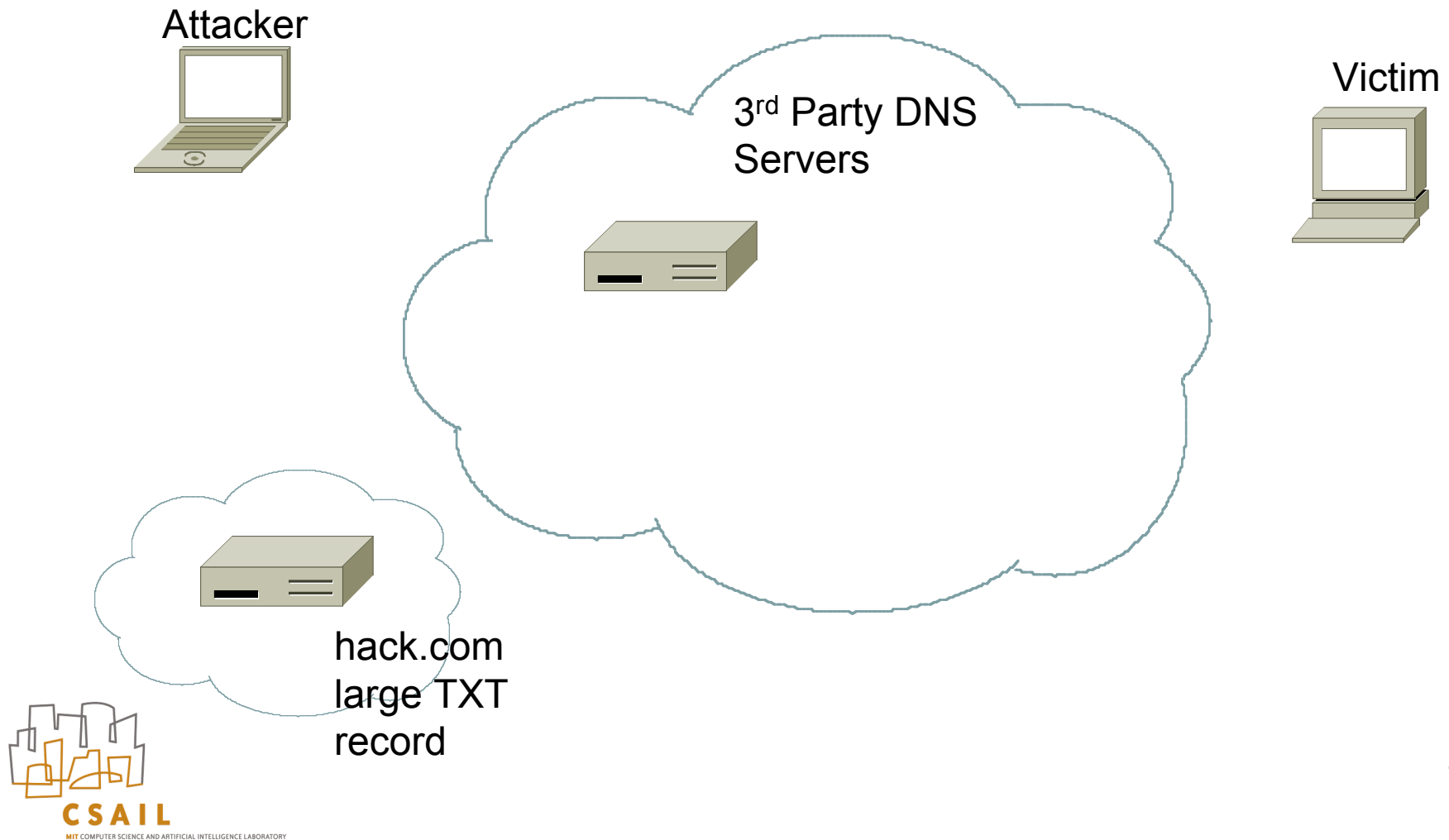
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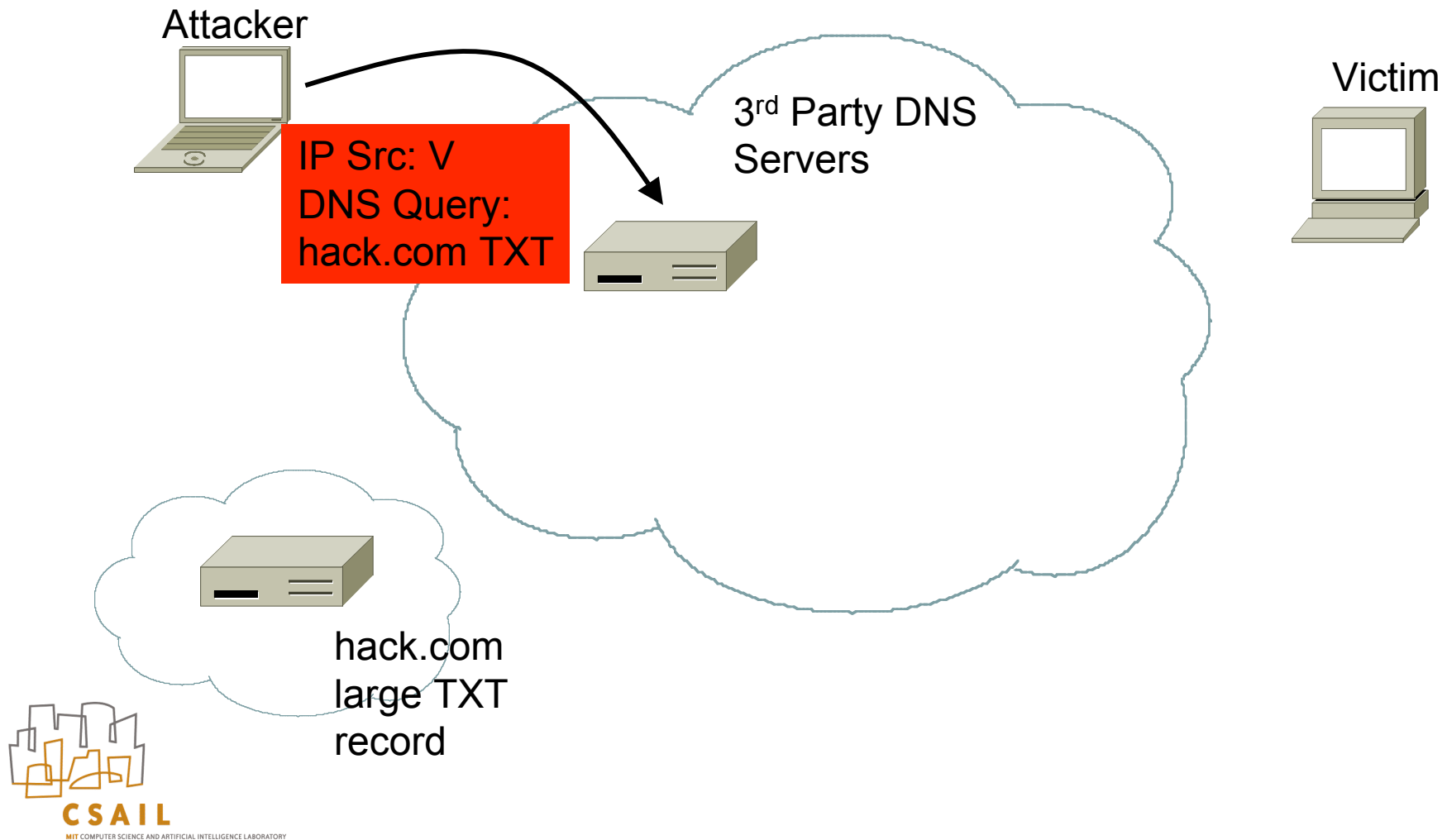
Prediction: spoofing increasingly a problem in the future

- **Spoofed traffic complicates a defenders job**
- Tracking spoofs is operationally difficult:
 - [Greene, Morrow, Gemberling NANOG 23]
 - Hash-based IP traceback [Snoeren01]
 - ICMP traceback [Bellare00]
- Consider **Slide from SRUTI 2005**
 - Today (in a ~~more case economy~~ more spoofing zombies are widely distributed, a network operator must defend against attack packets from 5% of routeable netblocks.
 - Future: if 25% of zombies capable of spoofing significant volume of the traffic could appear to come any part of the IPv4 address space
- Adaptive programs that make use of all local host capabilities to amplify their attacks

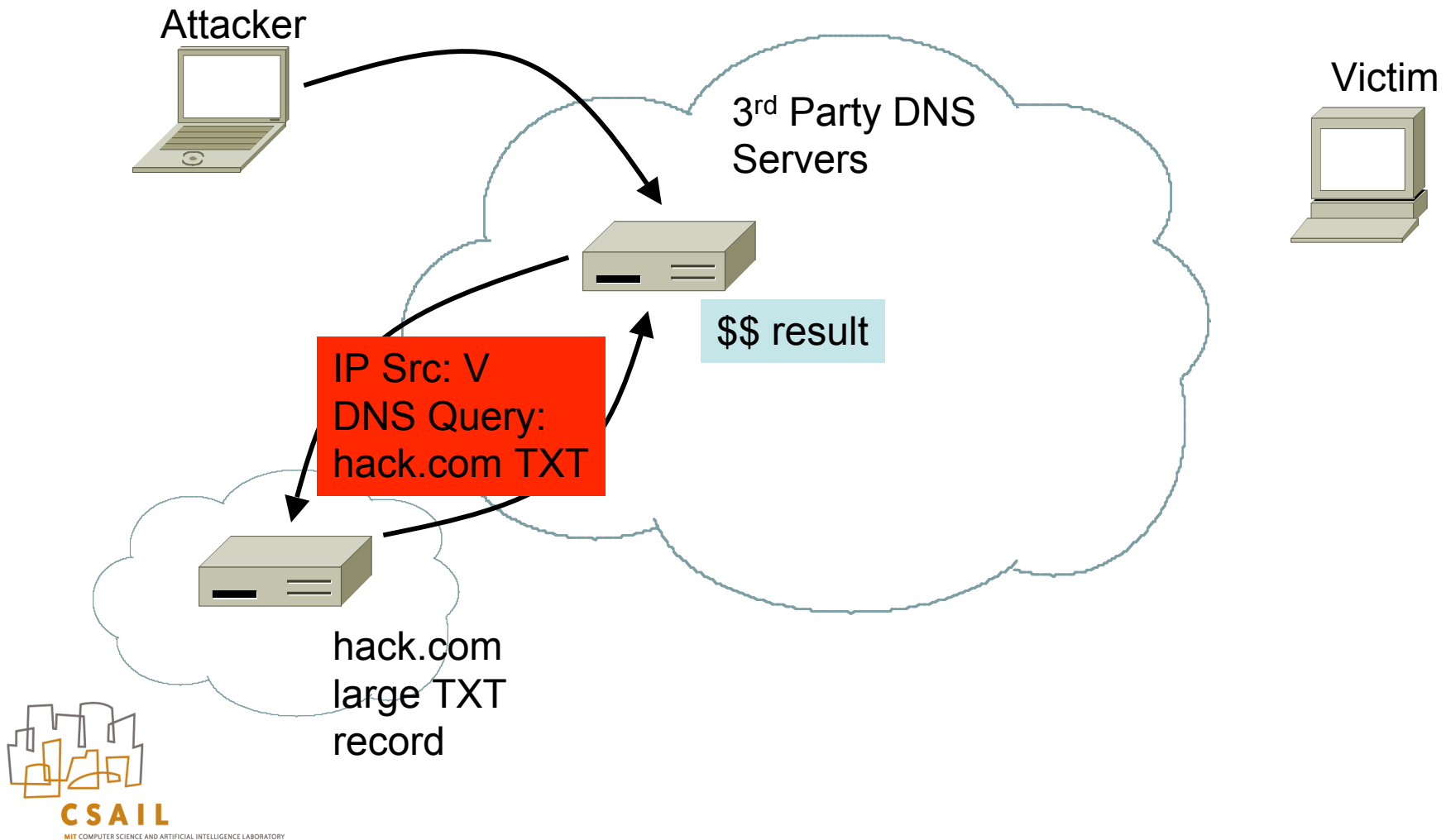
Prominent 2008 Example: DNS Amplifier Attack



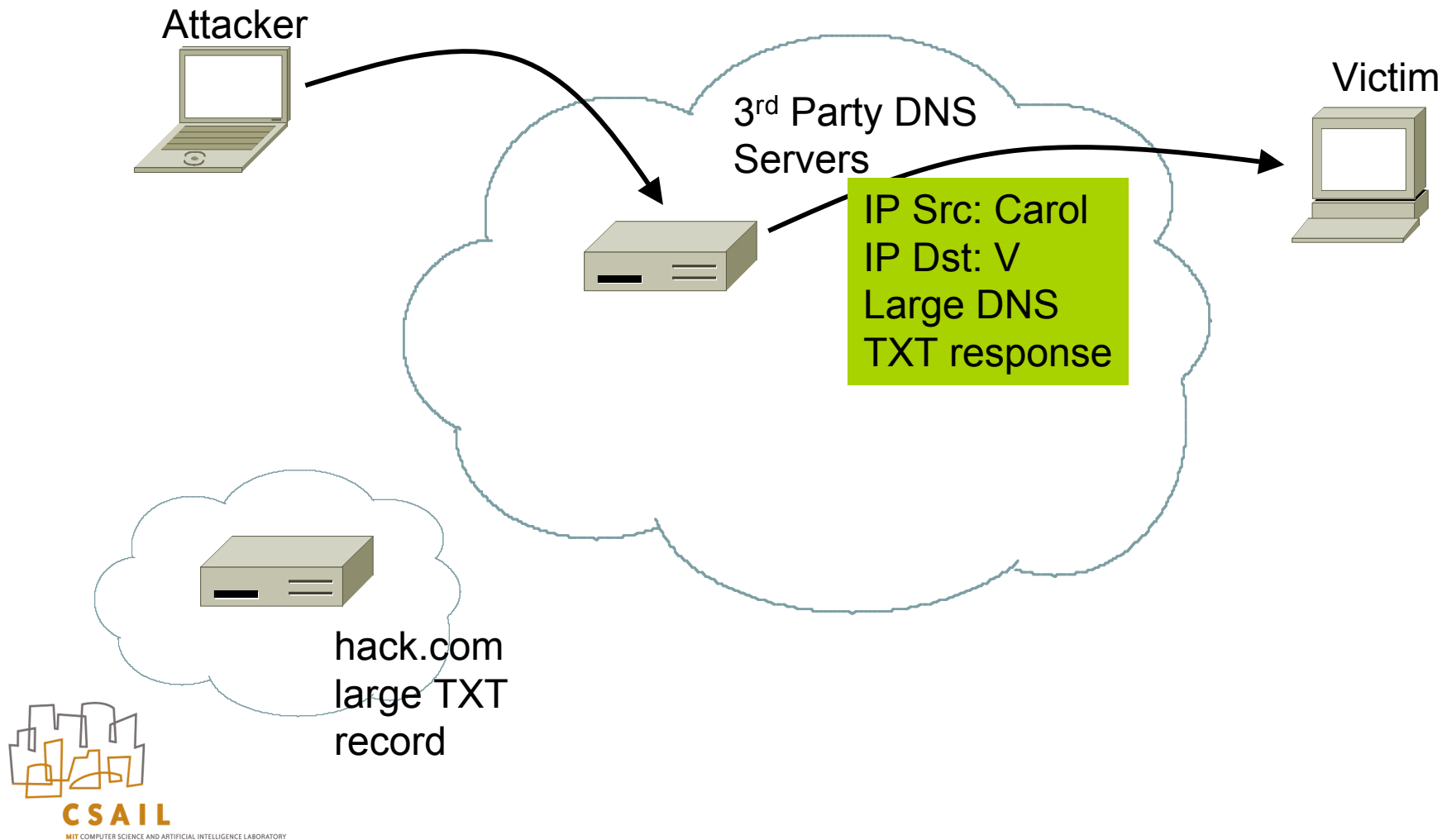
Prominent 2008 Example: DNS Amplifier Attack



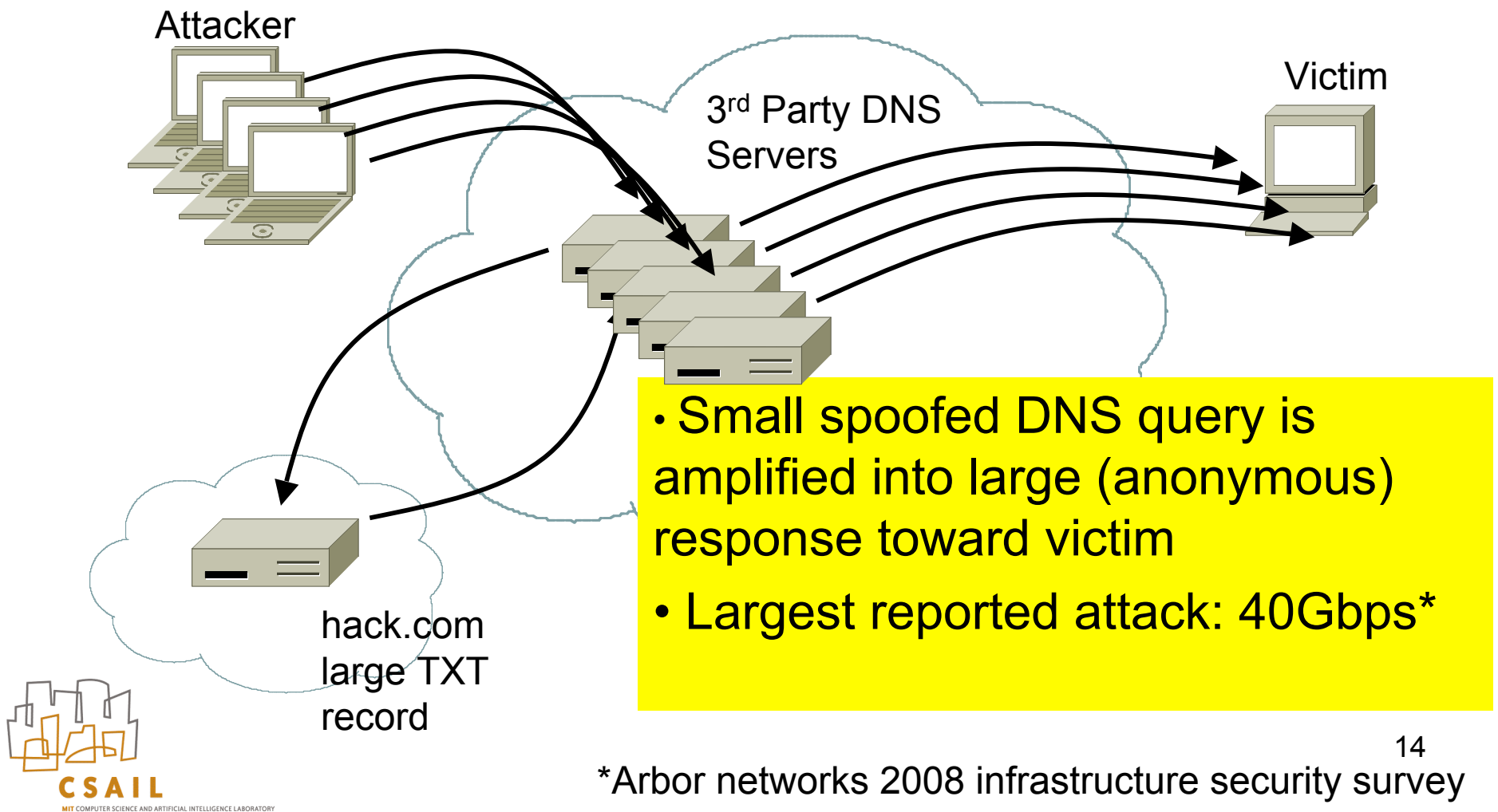
Prominent 2008 Example: DNS Amplifier Attack



Prominent 2008 Example: DNS Amplifier Attack



Prominent 2008 Example: DNS Amplifier Attack



Reasons to Believe Spoofing Matters (2009)

- DNS Amplifier Attacks
- In-Window TCP Reset Attacks
- Spam Filter Circumvention
- DNS Cache Poisoning
- UW reverse traceroute
- Spoofer web site statistics

The Operational Side

- Arbor:
 - “Reflective amplification attacks responsible for the largest attacks exploit IP spoofing”
 - “No bots were used in this attack. The attacker had a small number of compromised Linux boxes from which he’d launch the spoofed source DNS query.”
- What’s an operator to do?

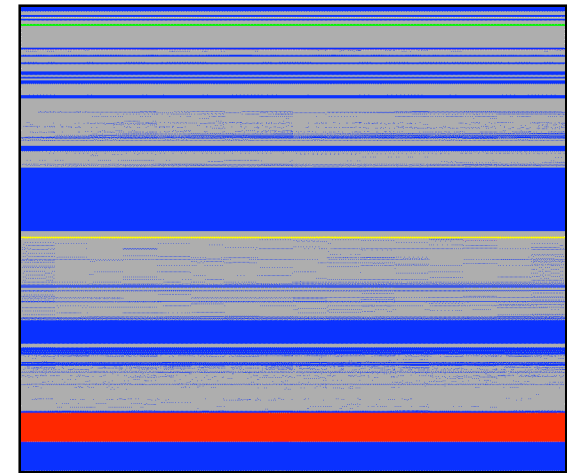
*Arbor networks 2008 infrastructure security survey

Operational View

- Not all sources are created equal
- IETF BCP38 best filtering practice

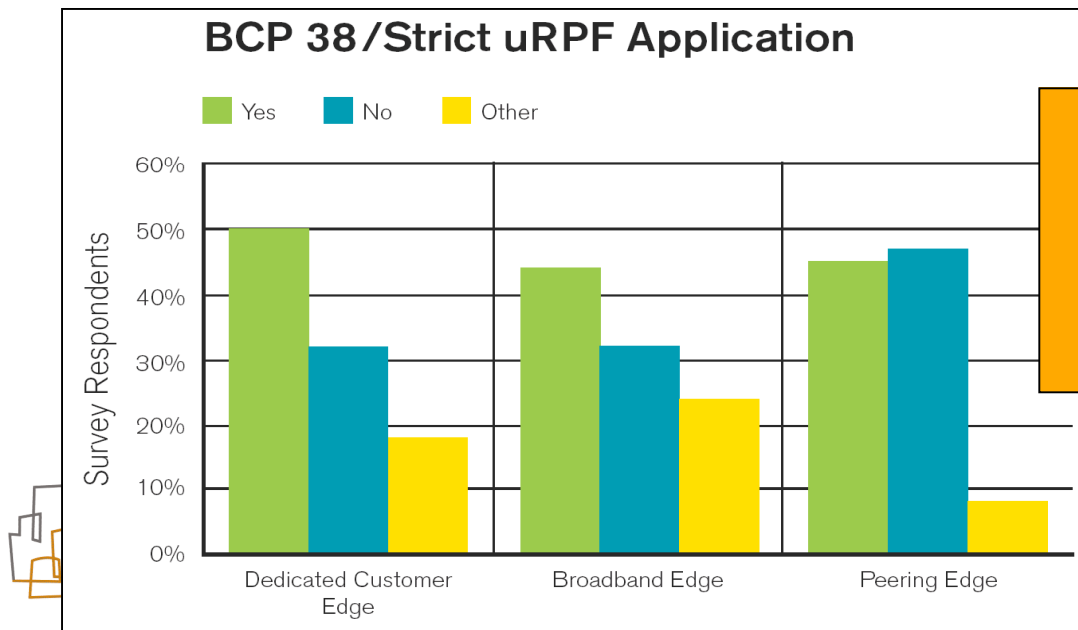
<u>Example Source IP</u>	<u>Description</u>	<u>Possible Defense</u>
1.2.3.4	Unallocated	Bogon Filters
6.1.2.3	Valid (In BGP table)	uRPF
192.168.1.1	RFC1918 private	Static ACL
Client IP ⊕ (2 ^N)	Neighbor Spoof	Switch, DOCSIS

IPv4 Address Space



Operational View

- We have defenses, what's the problem?
- BCP38 suffers from:
 - Incentive problem
 - Lack of hardware support (see NANOG)
 - Management nightmare (edge filters)



> 30% don't filter!
*Arbor networks 2008 infrastructure security survey

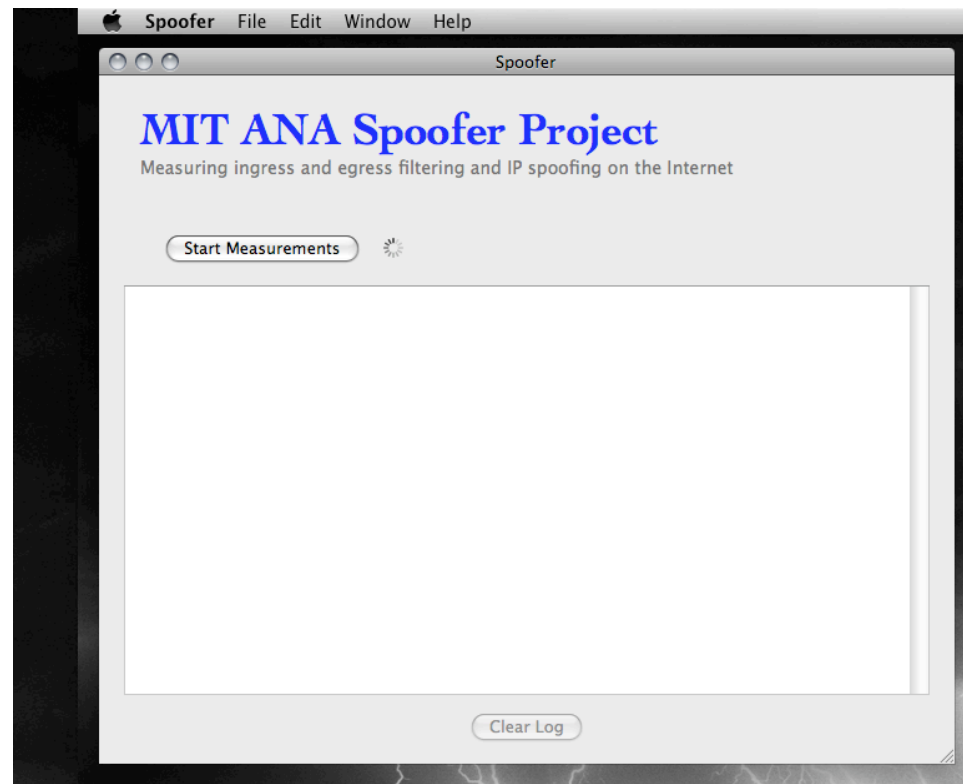
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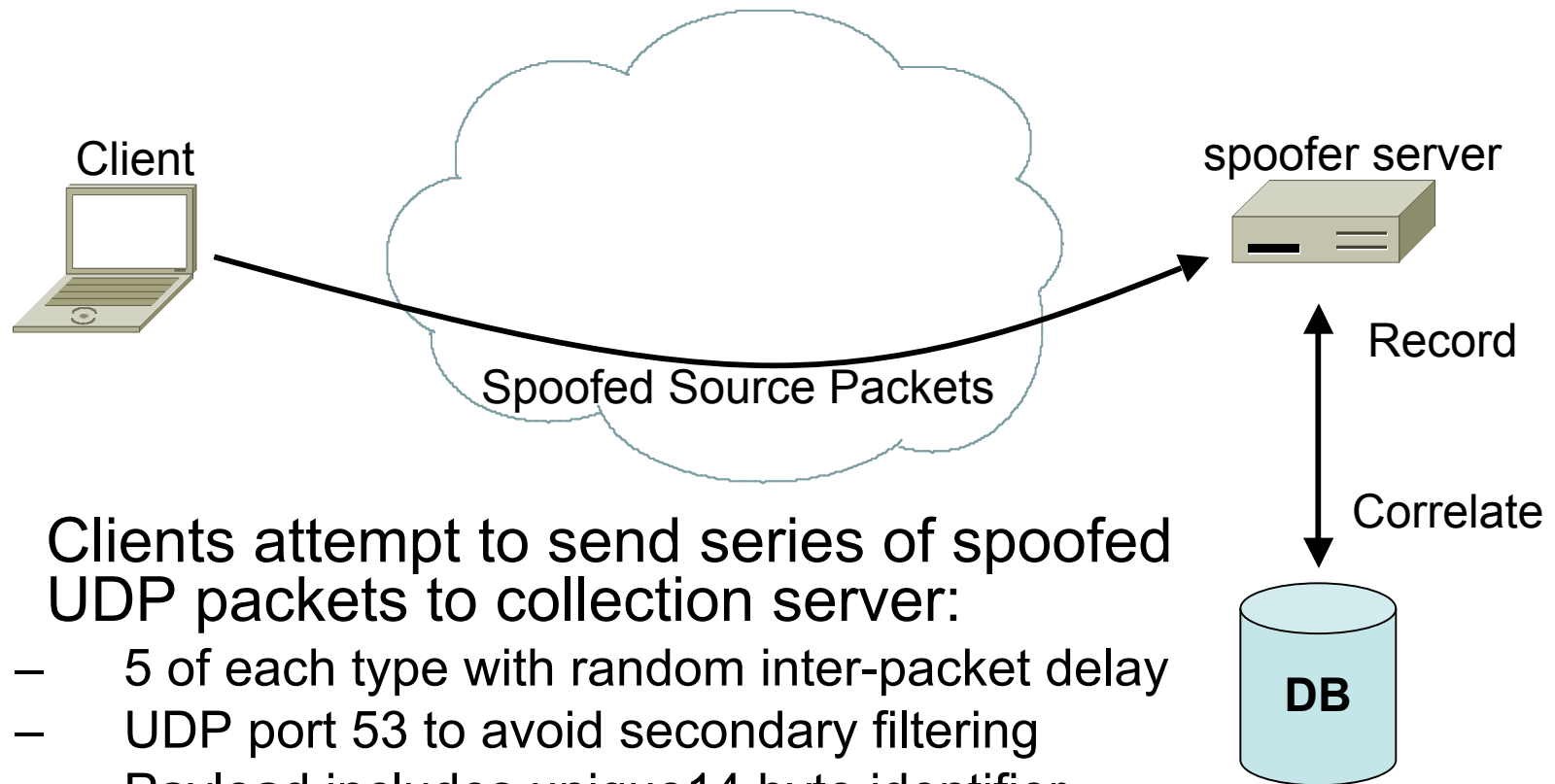
Spoofers Test Client



- Willing participants run “spoofers” client to test policy, perform inference, etc.
 - Binaries, source publicly available

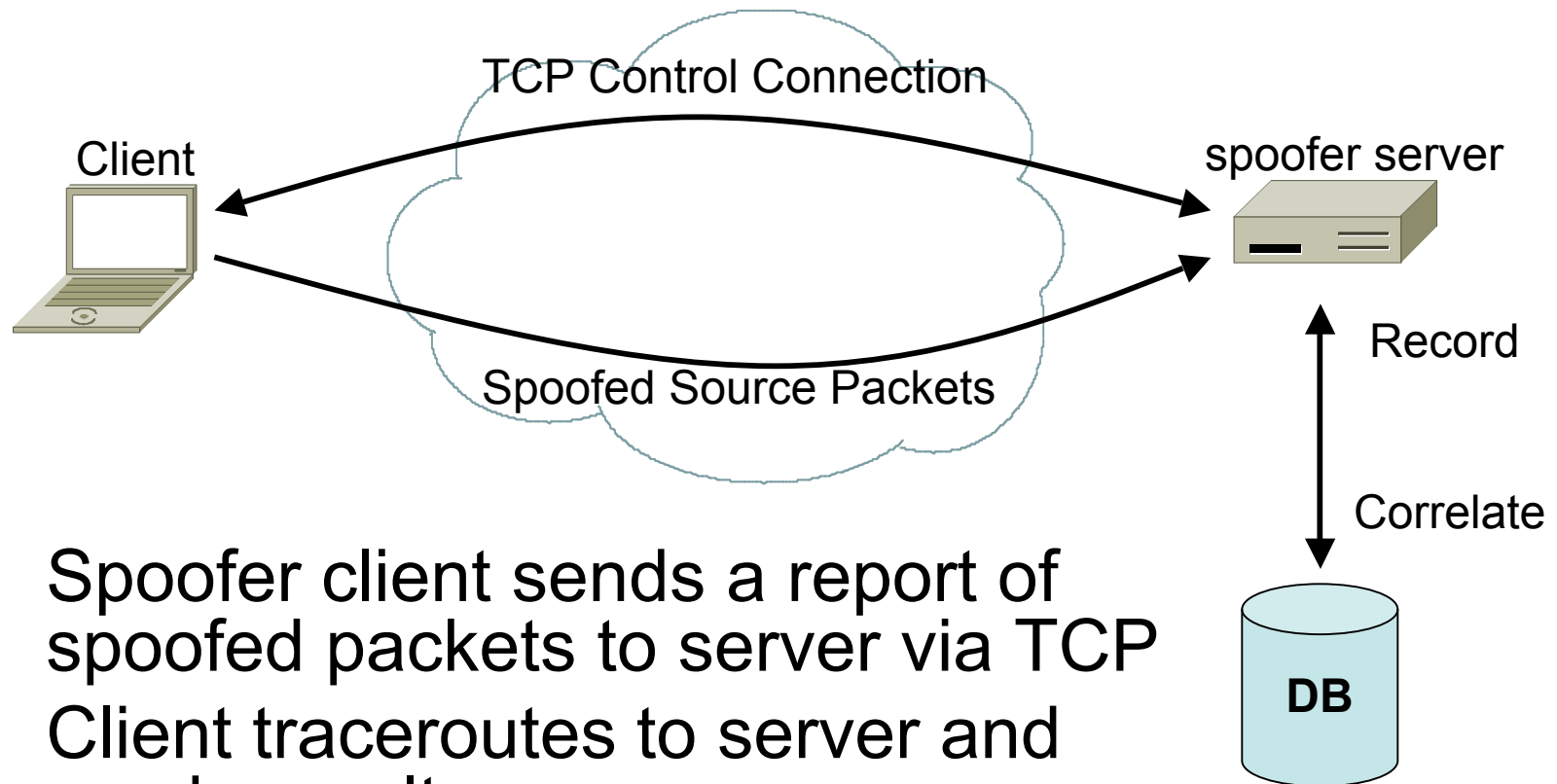


Spoofing Operation



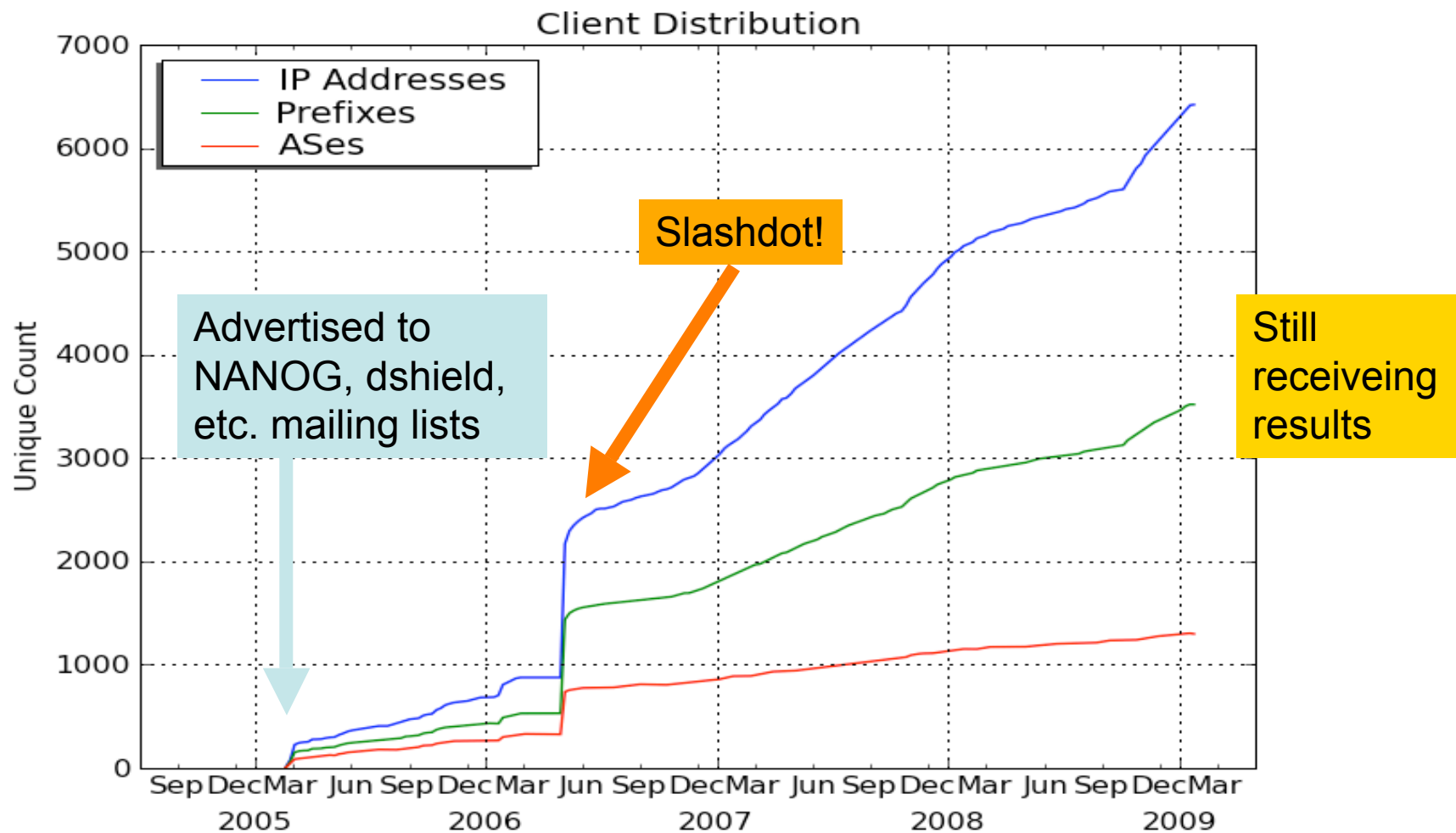
- Clients attempt to send series of spoofed UDP packets to collection server:
 - 5 of each type with random inter-packet delay
 - UDP port 53 to avoid secondary filtering
 - Payload includes unique 14 byte identifier
- Server stores received packets in DB

Spoofing Operation

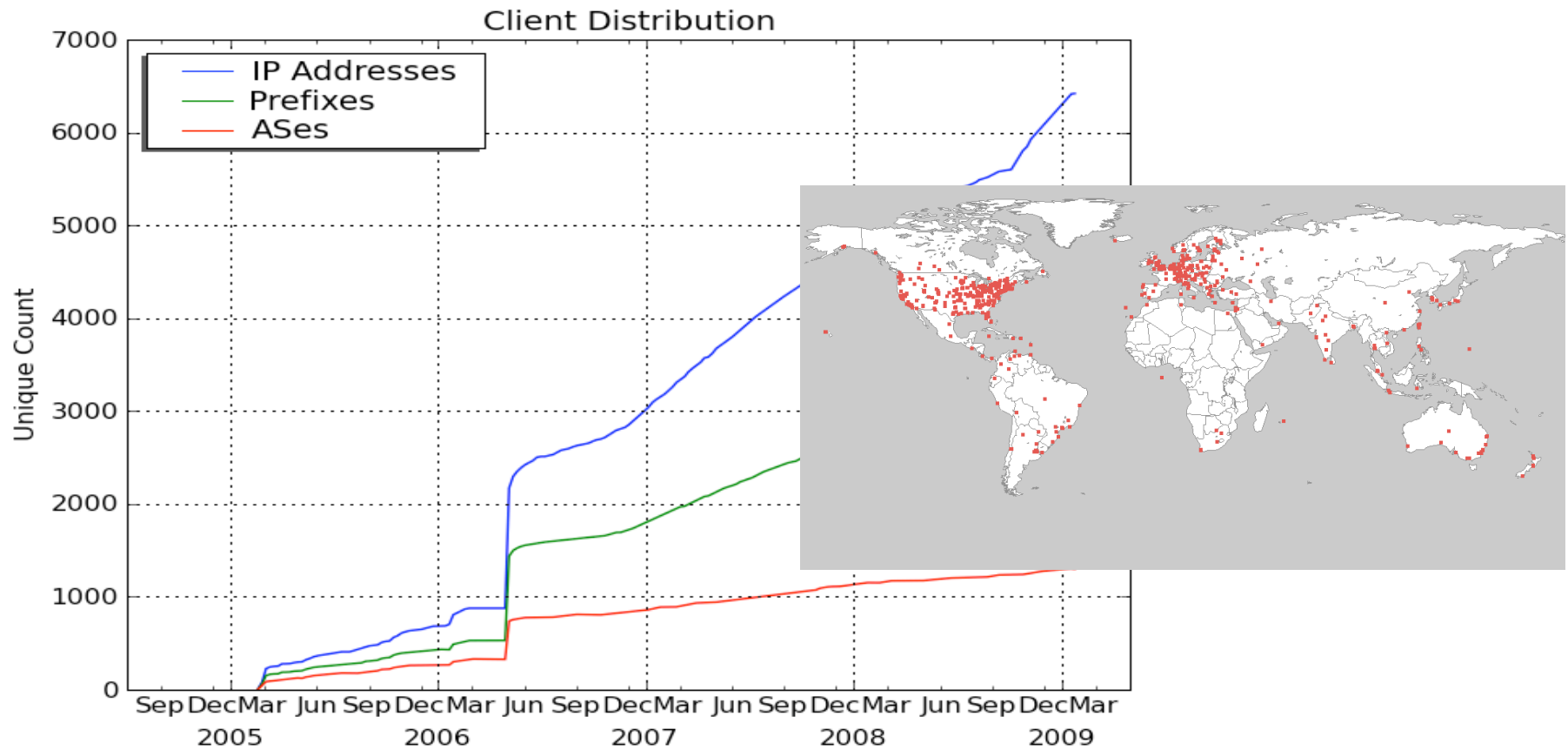


- Spoofer client sends a report of spoofed packets to server via TCP
- Client traceroutes to server and sends result
- TCP destination port 80 used to avoid secondary filtering effects

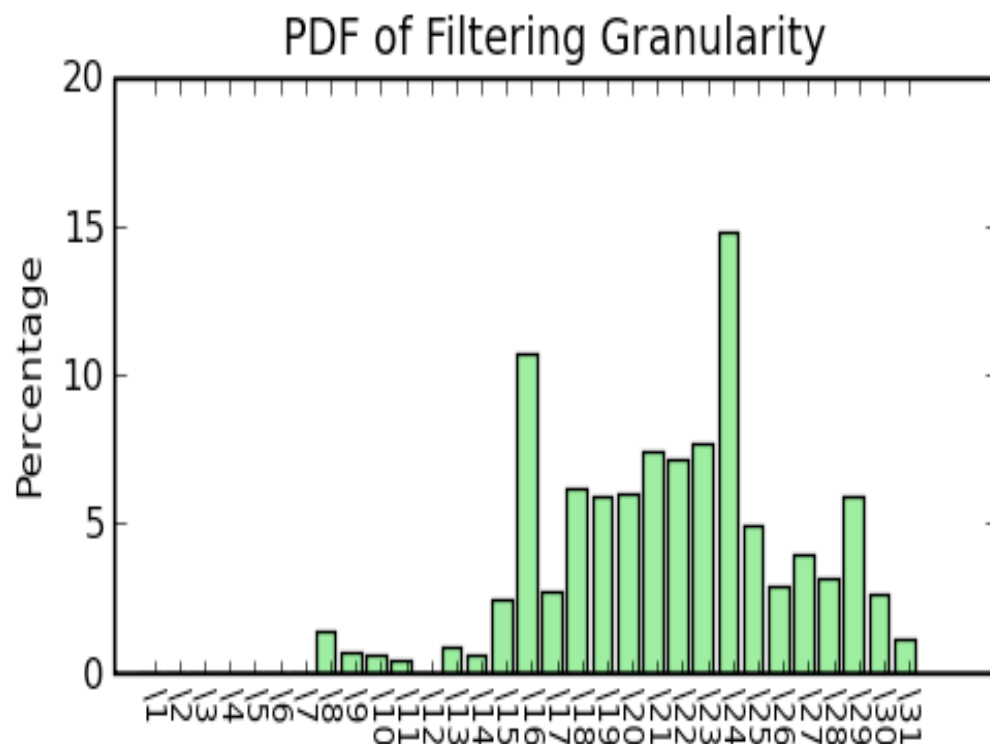
Client Population



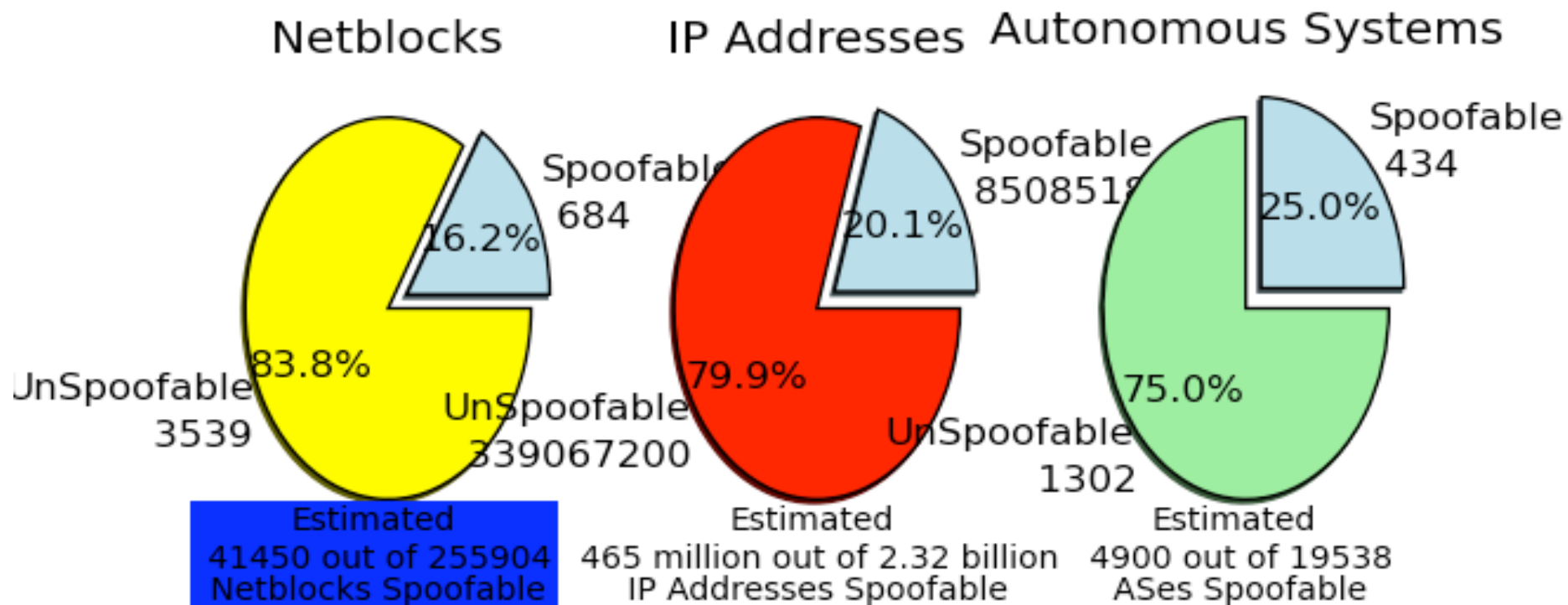
Client Population Distribution



Filtering Specificity



- Clients test own IP \oplus (2^n) for $0 < n < 24$
- Filtering on a /8 boundary enables a client within that network to spoof ~16M addresses
- >30% of clients “unable” to spoof can spoof neighbors
- Exclude “neighbor spoof” from macro results



- **Spoofable:** spoofing of private, unallocated, or valid IP packets possible from these locations
- Agrees to a first-order with Arbor survey
- But... these numbers cause even more disagreement!

Spoofers Project

- Background
- Recent Relevance
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- **What's New: Methodology**
- What's New: Data
- Parting Thoughts

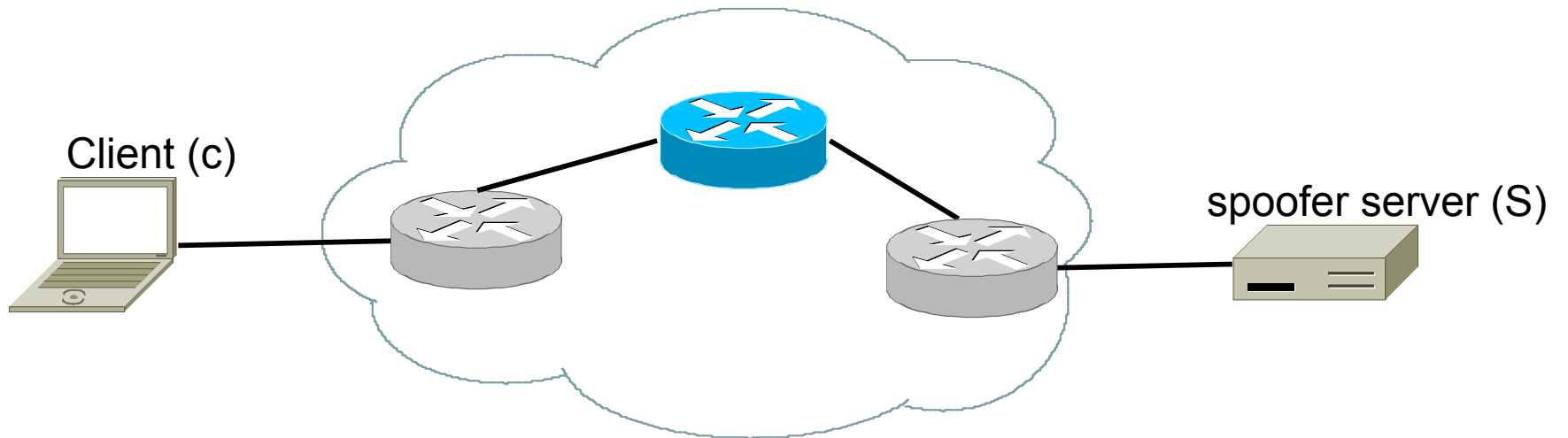
What's New: Methodology

- Goal:
 - Resolve ambiguity
 - Increase confidence
- New:
 - tracefilter
 - Tied into CAIDA's ark distributed measurement infrastructure
 - More detailed analysis
 - Longitudinal analysis over four-years of data

tracefilter

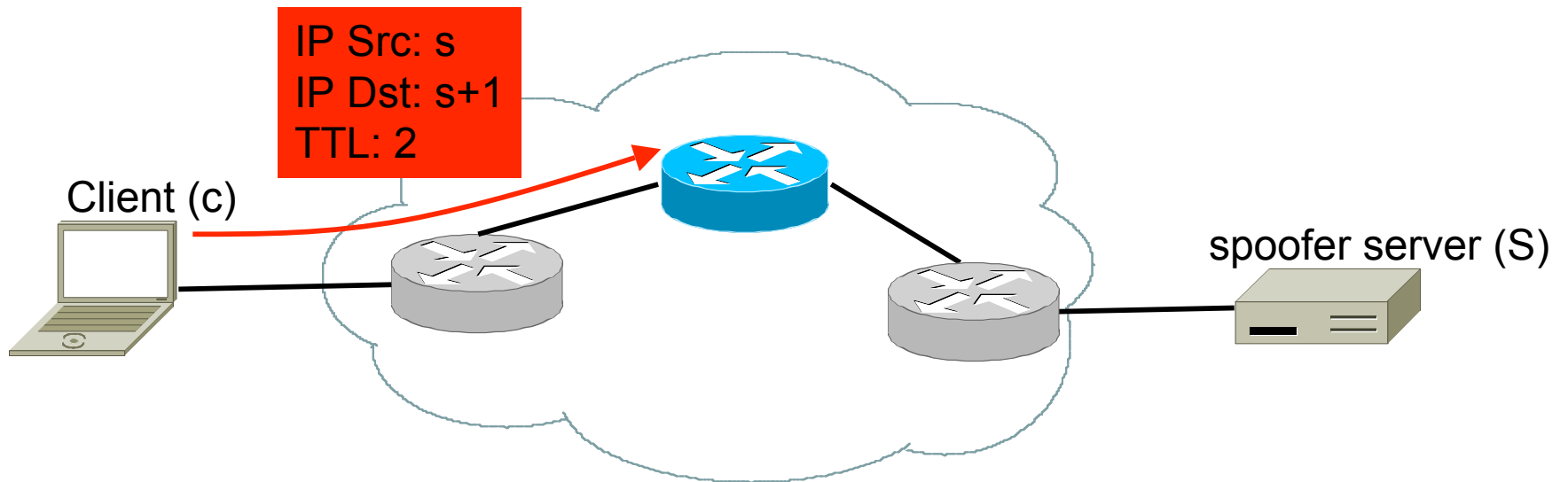
- A tool for *locating* source address validation (anti-spoofing) filters along path
- “traceroute for BCP38”
- Better understand who is/is not filtering

tracefilter



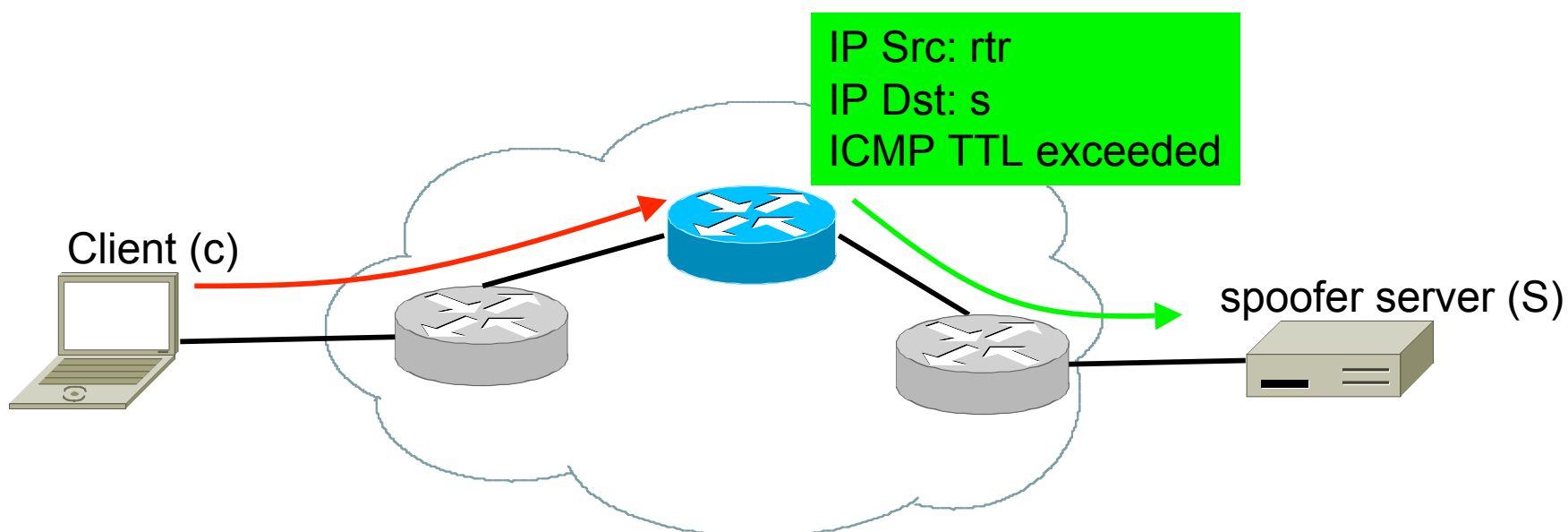
- Client c works in conjunction with our server S

tracefilter



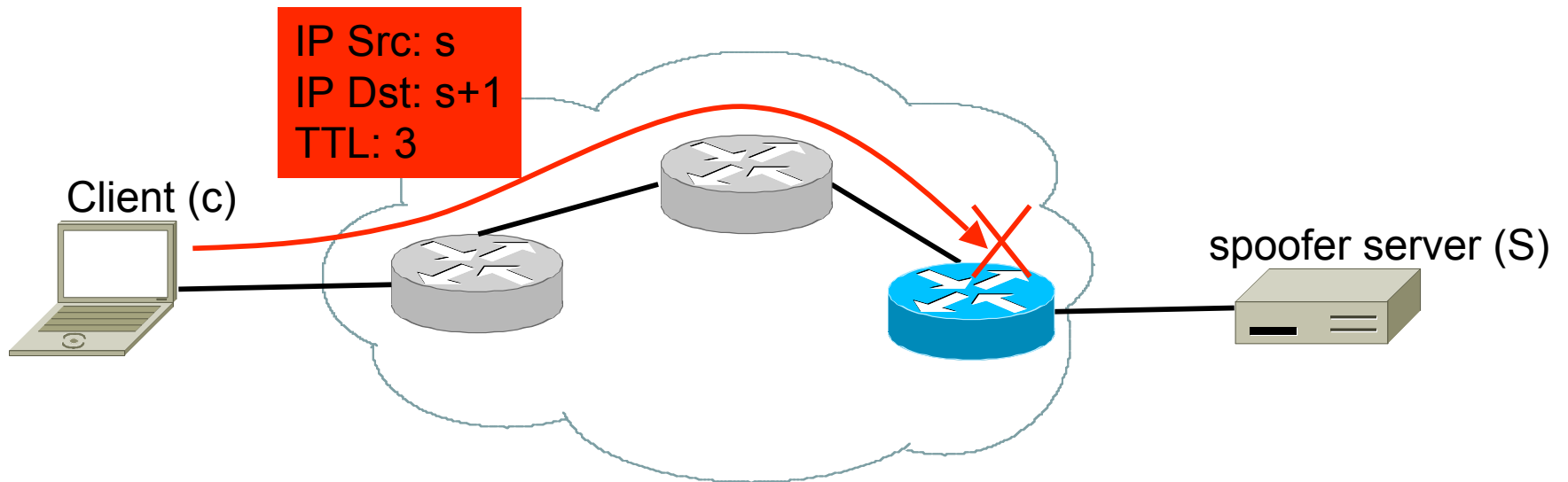
- c sends spoofed packet with:
- $t\text{tl}=x$, $\text{src}=\text{S}$, $\text{dst}=\text{S}+1$ for $0 < x < \text{pathlen}$

tracefilter



- S receives ICMP expiration messages from routers along path
- For each decoded TTL, S records which spoofed packets are received

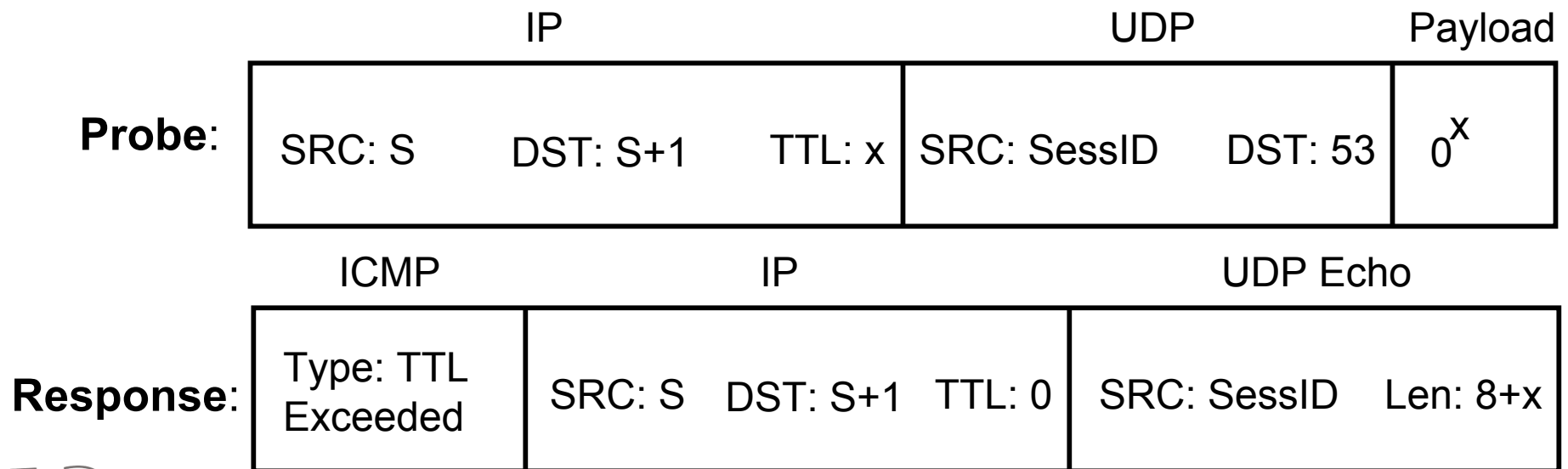
tracefilter



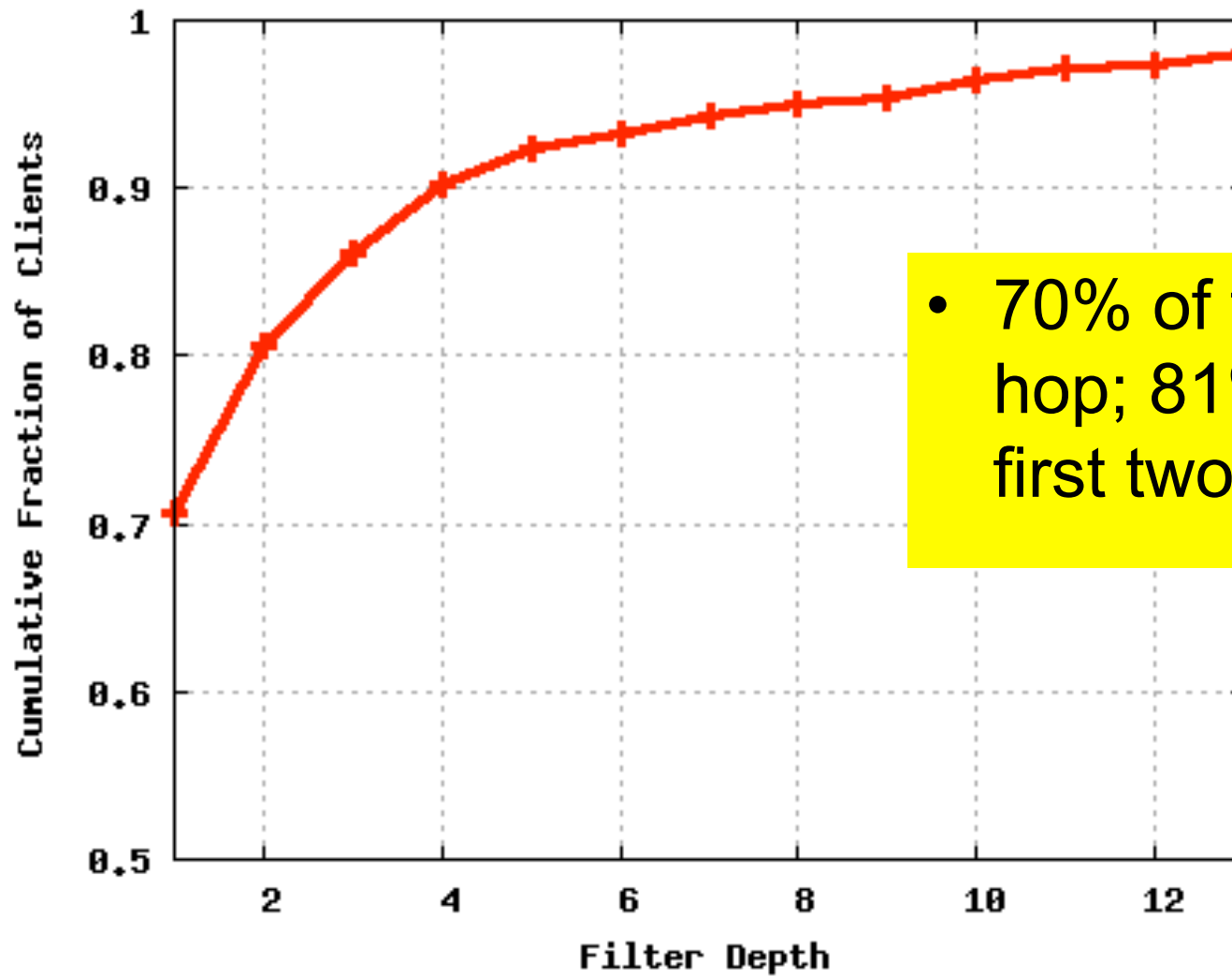
- Increase TTL, repeat
- Largest TTL indicates filtering point

tracefilter

- How can *S* determine *originating* TTL of *c*'s packets?
- ICMP echo includes only 28 bytes of expired packet
- *c* encodes TTL by padding payload with zeros

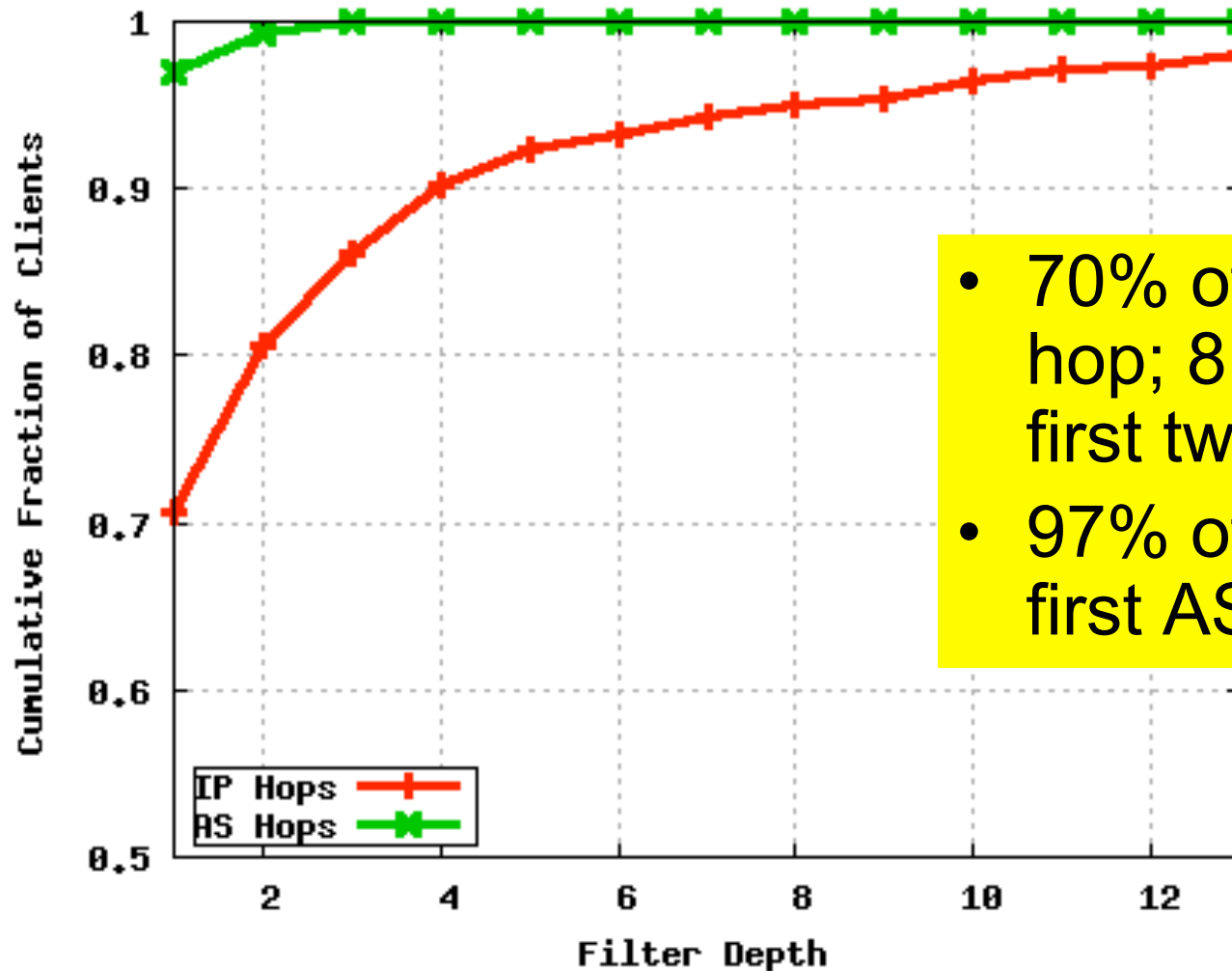


tracefilter Results



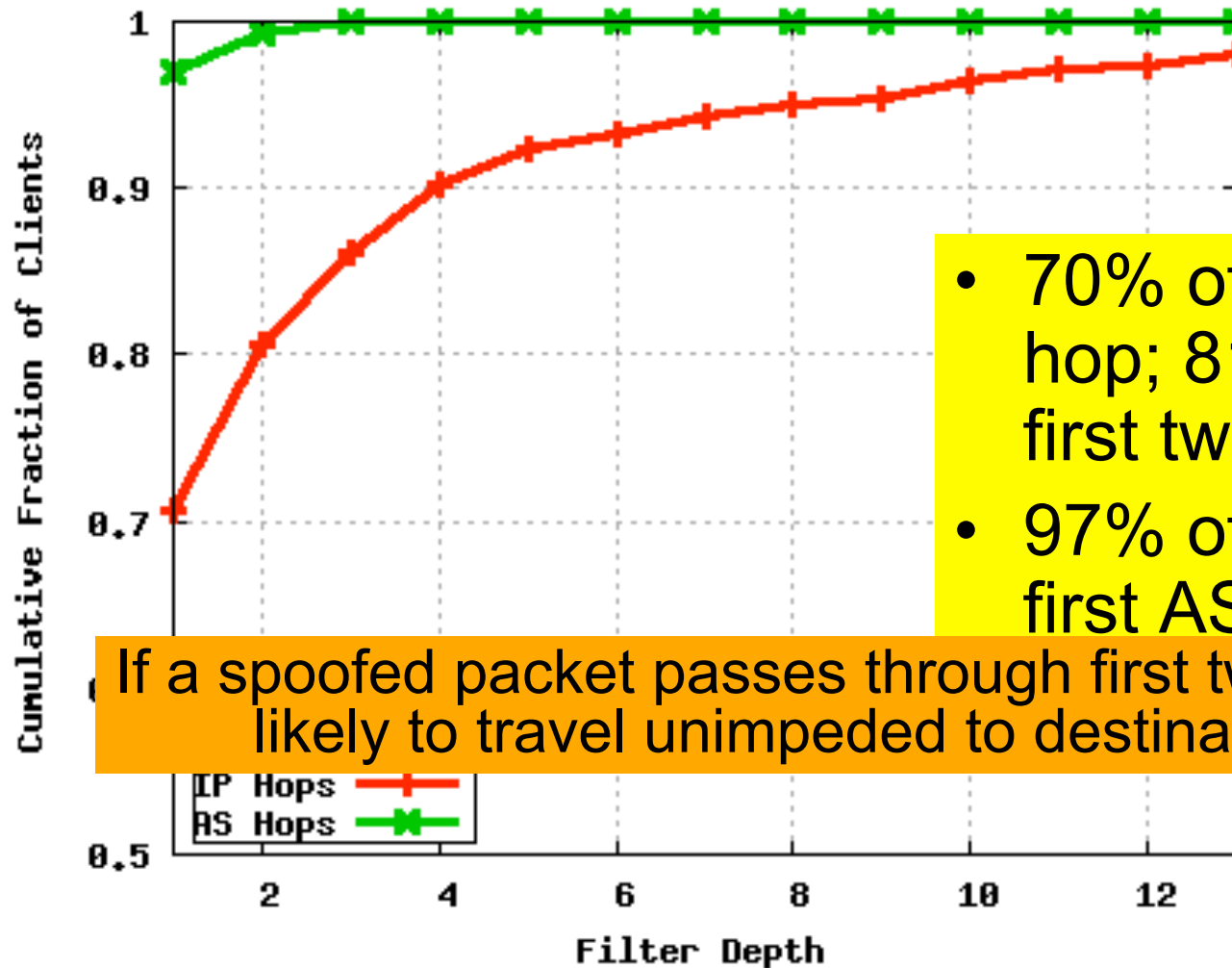
- 70% of filters at 1st hop; 81% within first two hops

tracefilter Results



- 70% of filters at 1st hop; 81% within first two hops
- 97% of filters within first AS

tracefilter Results



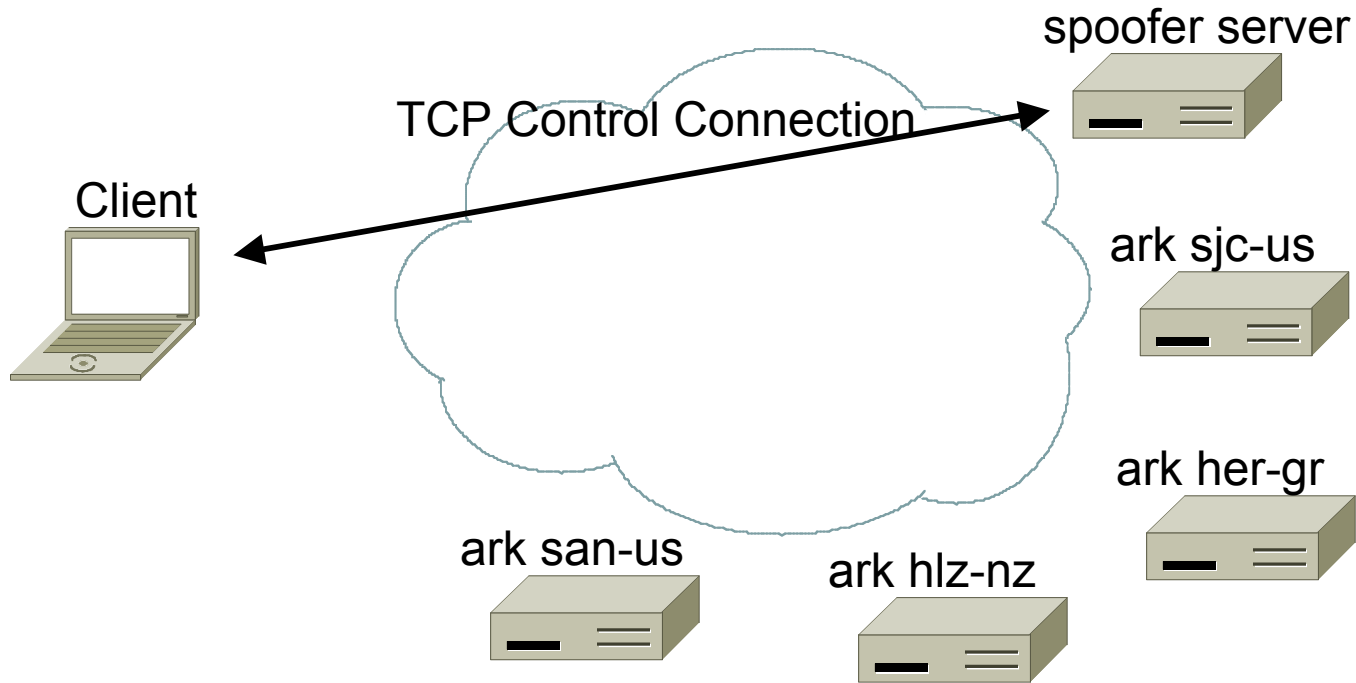
- 70% of filters at 1st hop; 81% within first two hops
- 97% of filters within first AS

If a spoofed packet passes through first two hops, likely to travel unimpeded to destination

Ark Support

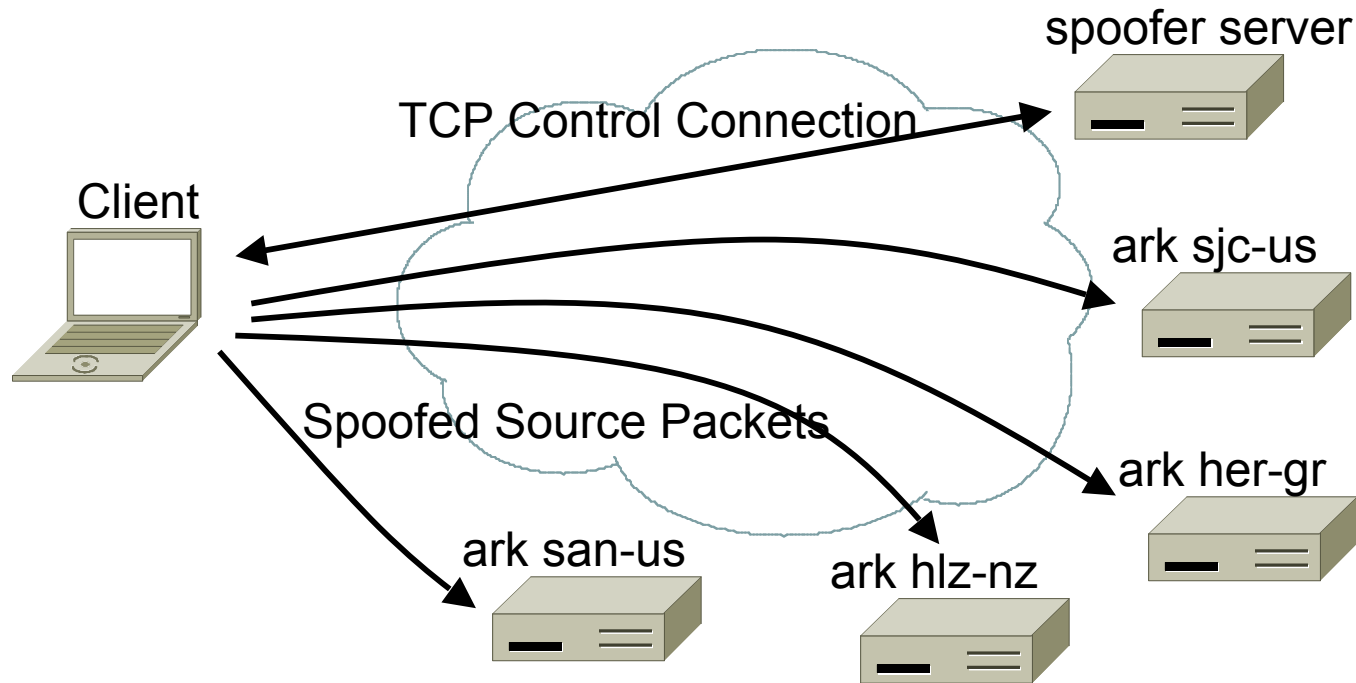
- Spoofer tester now tied into CAIDA's archipelago distributed measurement infrastructure (Ark)
- Provides invaluable additional inference capability
- Allows us to resolve aforementioned ambiguity

Utilizing Ark Infrastructure



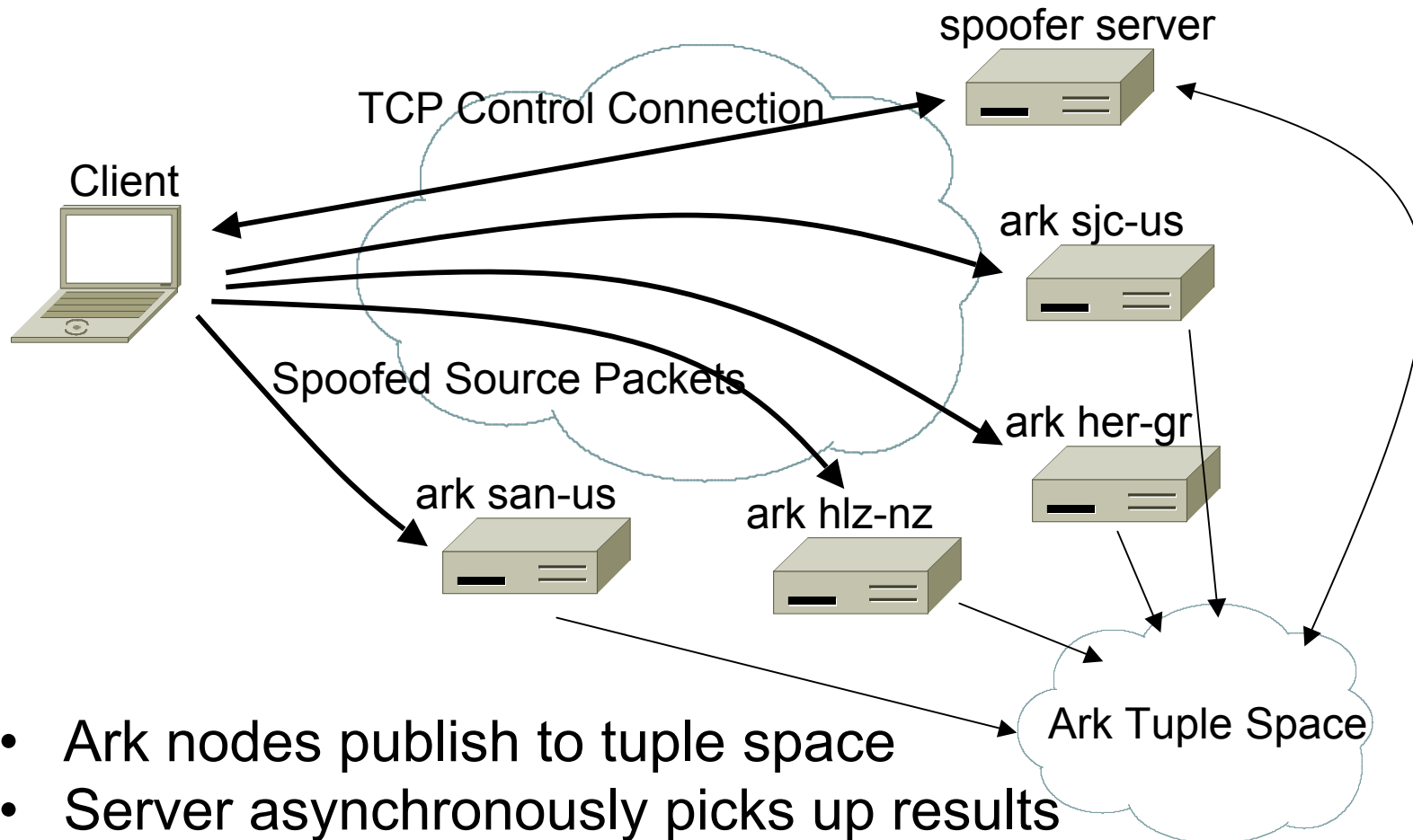
- Server and Ark nodes agree on common HMAC key
- Provide client with (SRC, DST, KEY, SEQ) tuples

Utilizing Ark Infrastructure



- Client sends HMAC keyed spoof probes to ark nodes
- Client runs traceroute to each ark node in parallel

Utilizing Ark Infrastructure

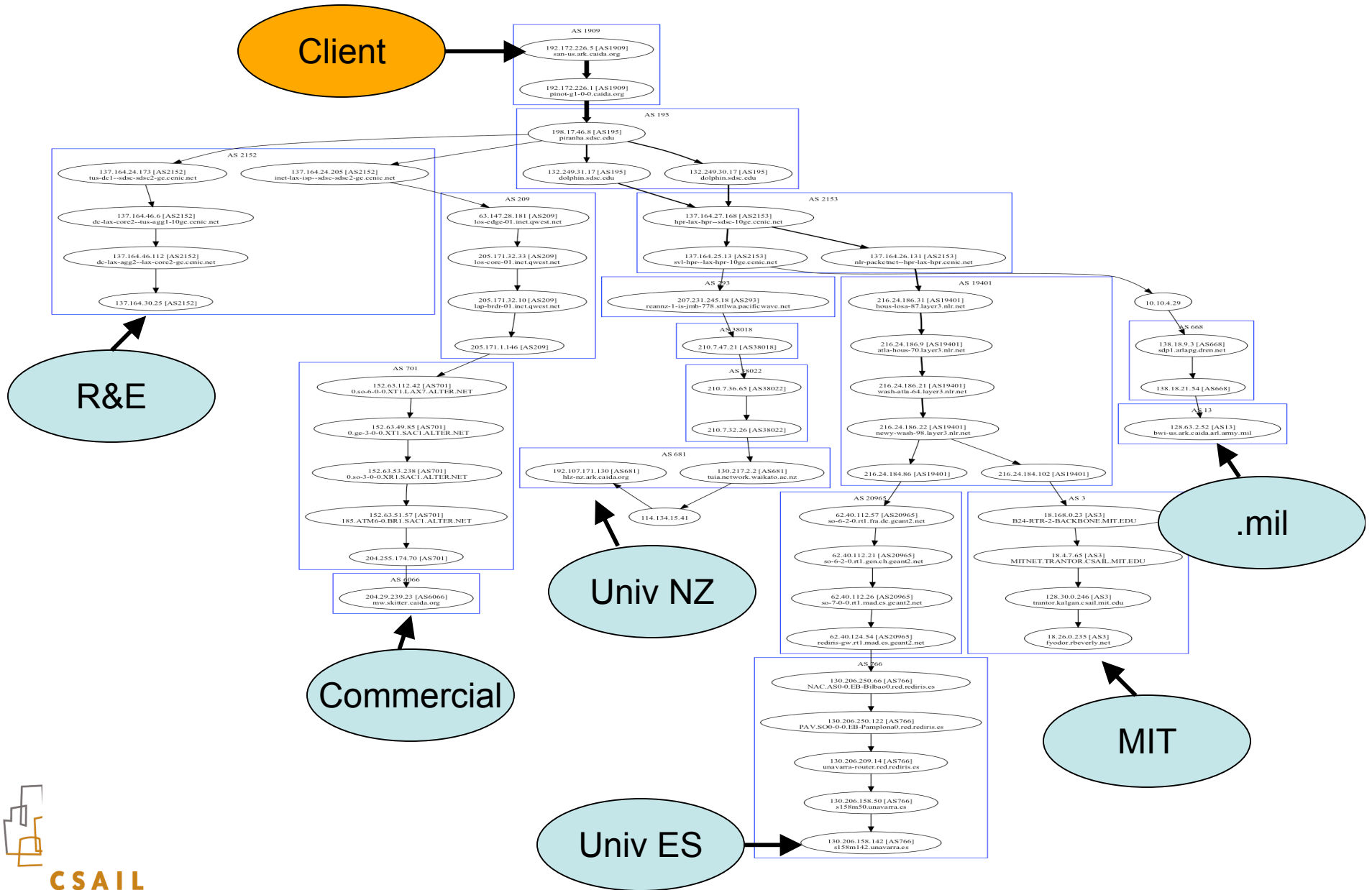


- Ark nodes publish to tuple space
- Server asynchronously picks up results

Value of Ark

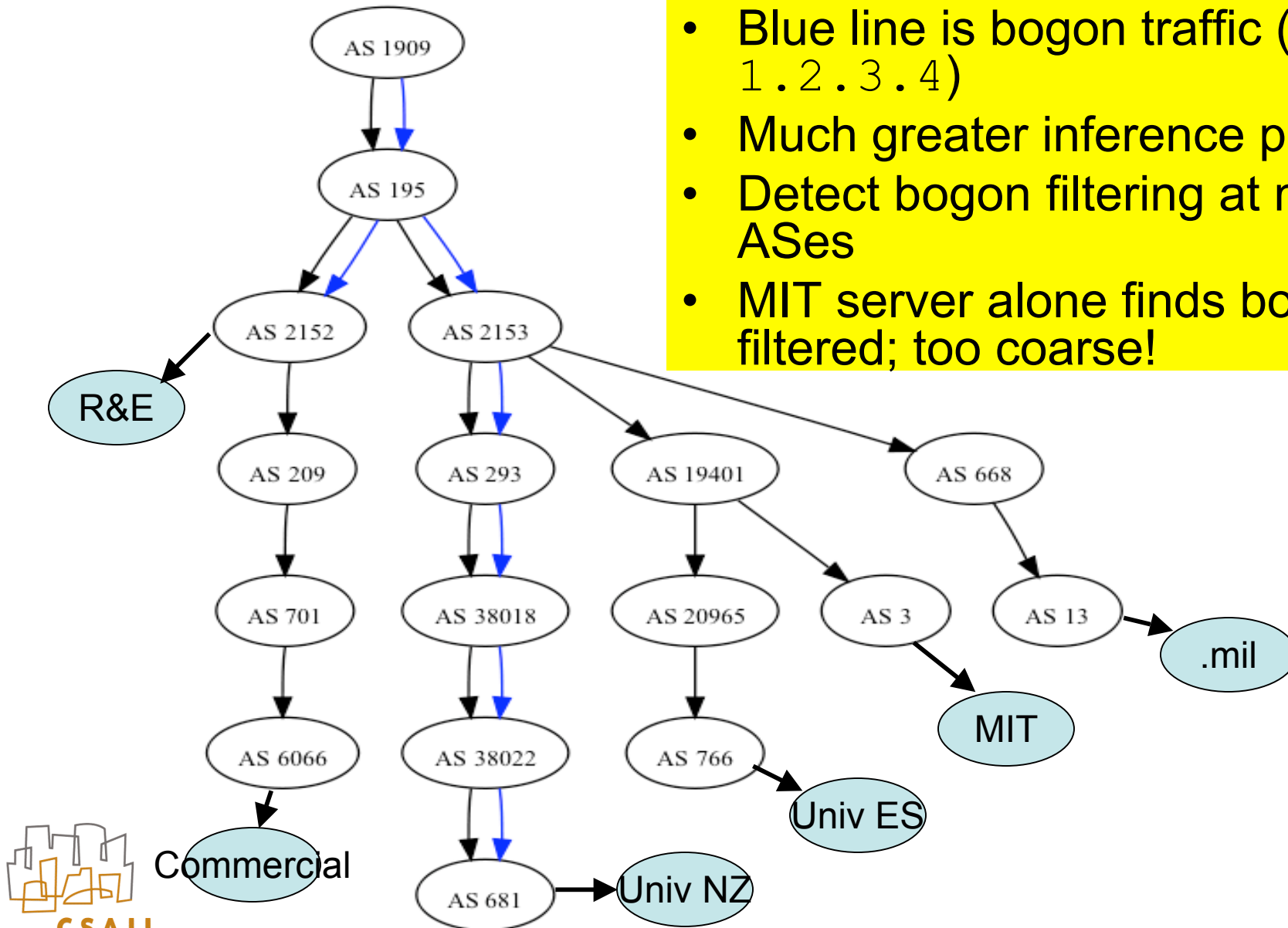
- How does Ark allow us better inference
- Example:

Multiple Destinations



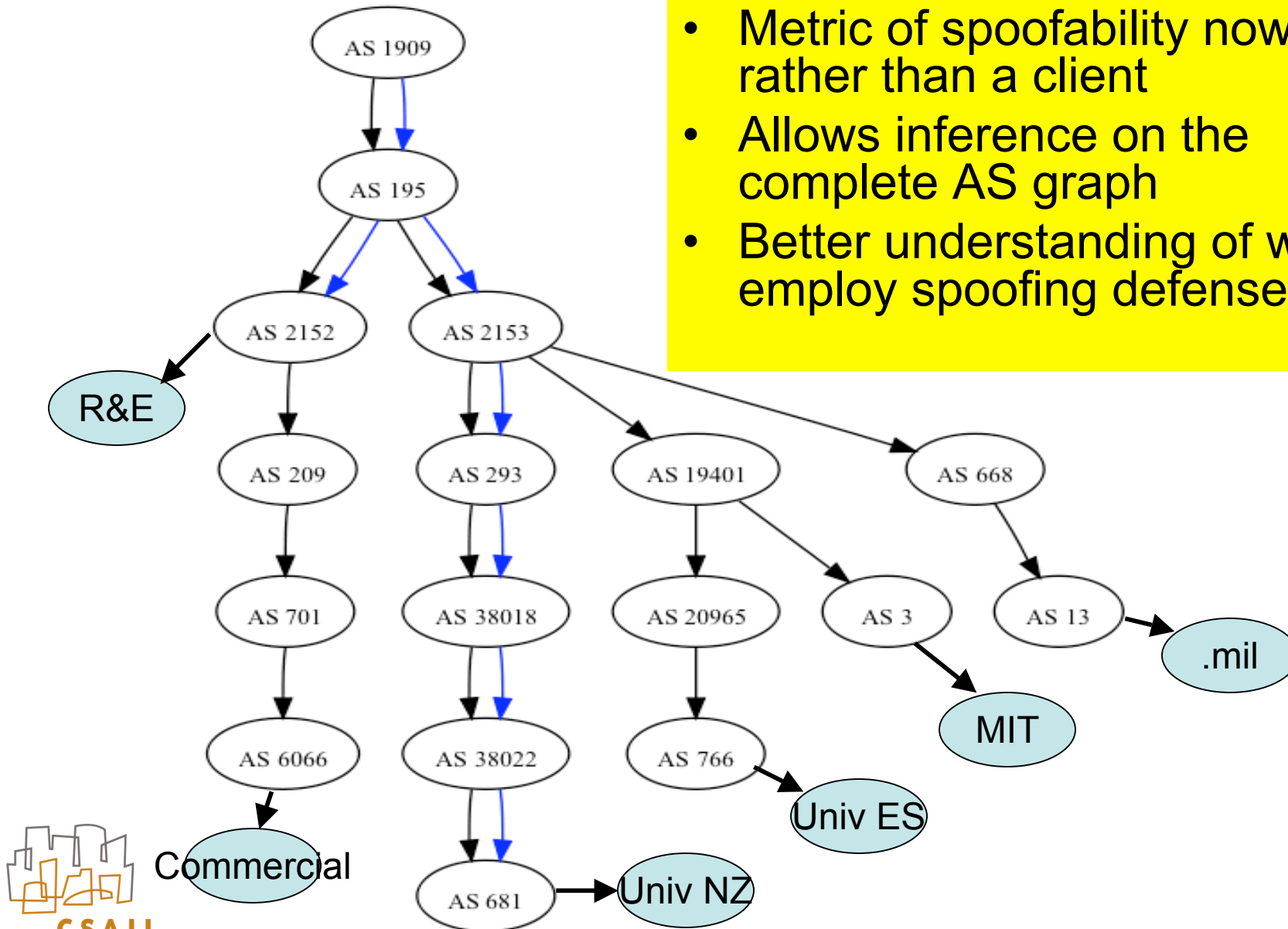
Multiple Destinations

- Blue line is bogon traffic (IP: 1.2.3.4)
- Much greater inference power
- Detect bogon filtering at multiple ASes
- MIT server alone finds bogons filtered; too coarse!



Multiple Destinations

- Metric of spoofability now a path rather than a client
- Allows inference on the complete AS graph
- Better understanding of where to employ spoofing defenses



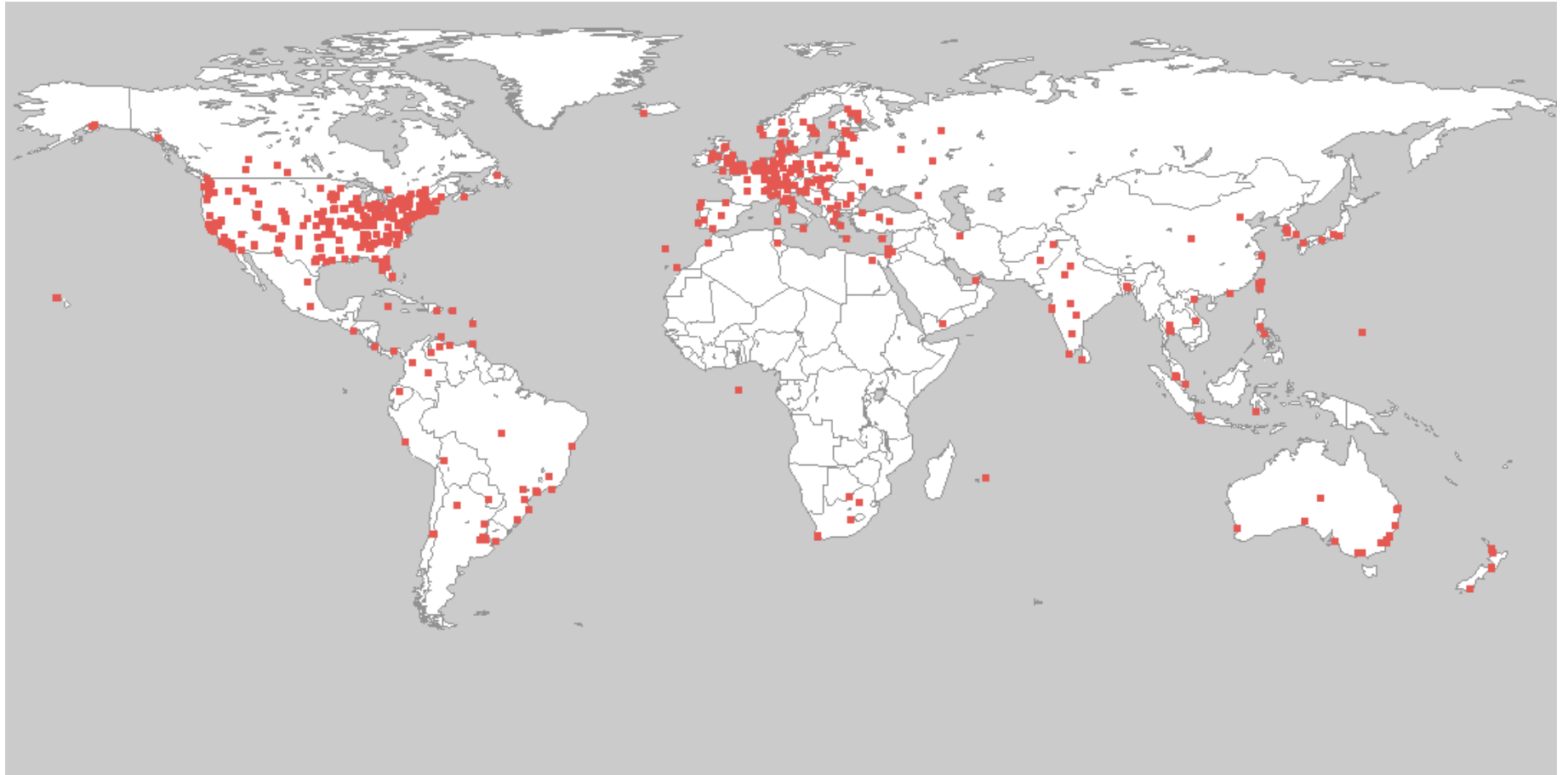
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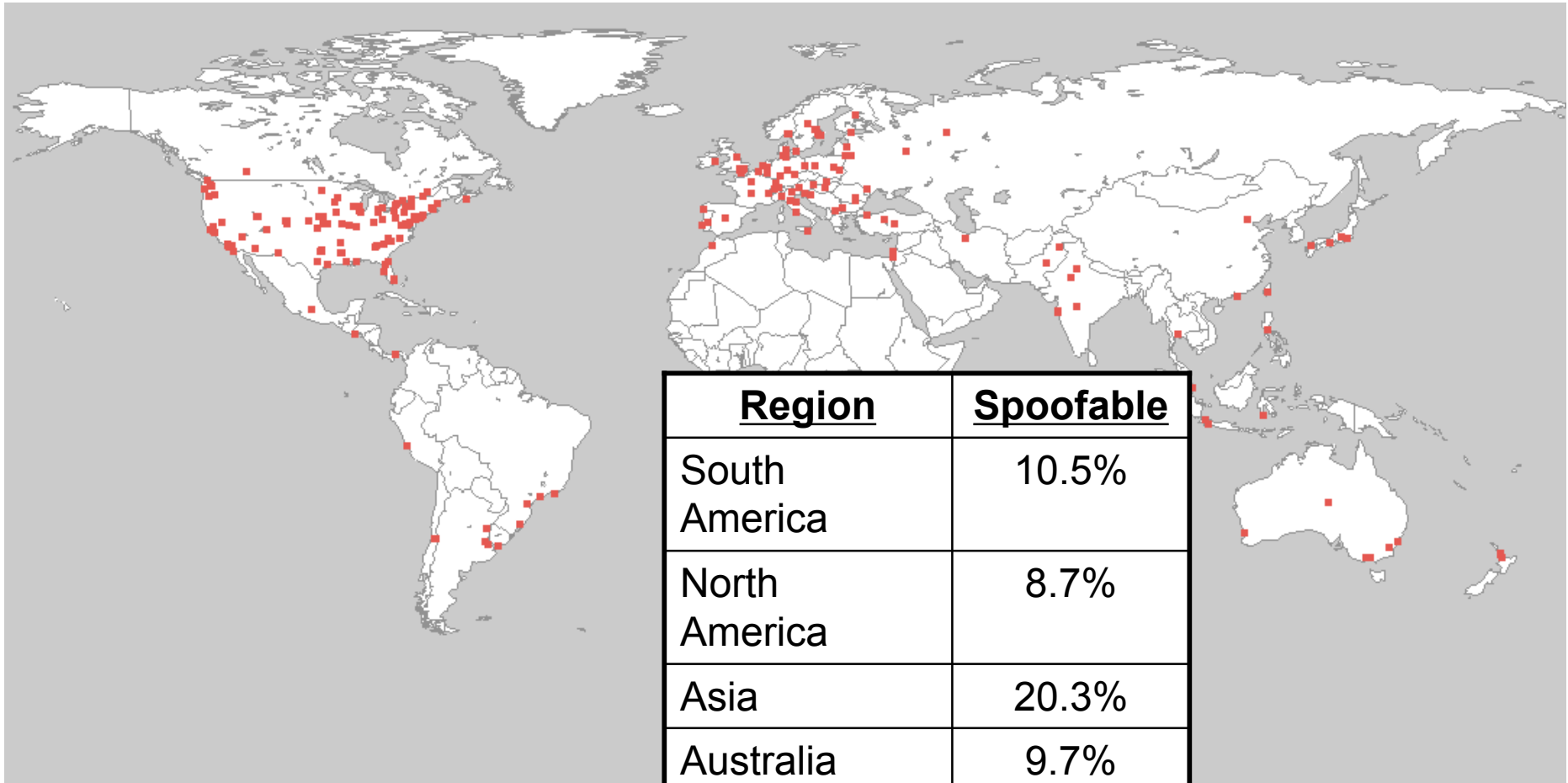
Deeper Analysis

- Question we want to answer:
 - Geographic analysis
 - Large or small providers filter?
 - What kinds of providers?

Geographic (Tests)

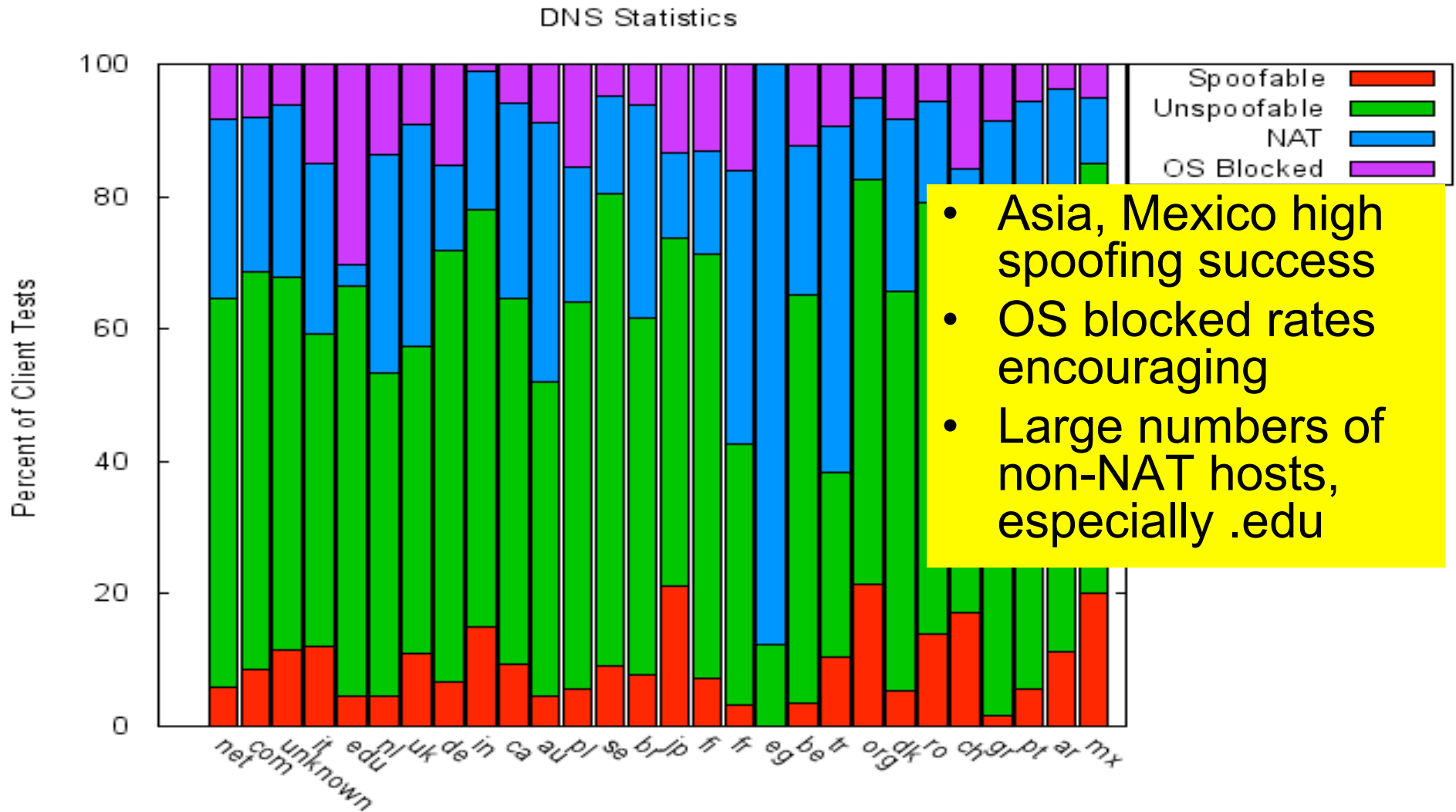


Geographic (Spoofable)



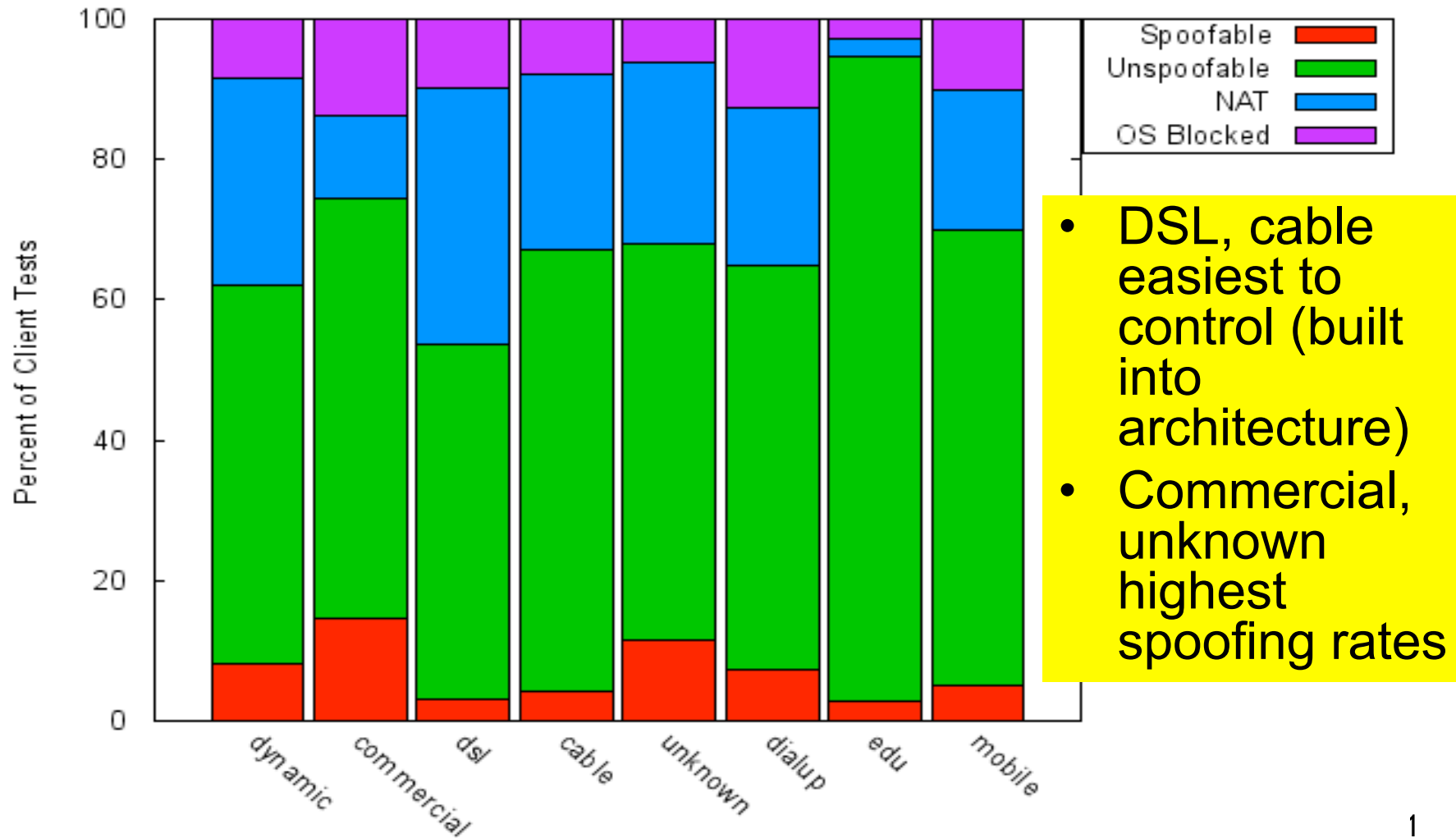
Region	Spoofable
South America	10.5%
North America	8.7%
Asia	20.3%
Australia	9.7%
Europe	9.2%
Africa	10.5%

DNS Stats



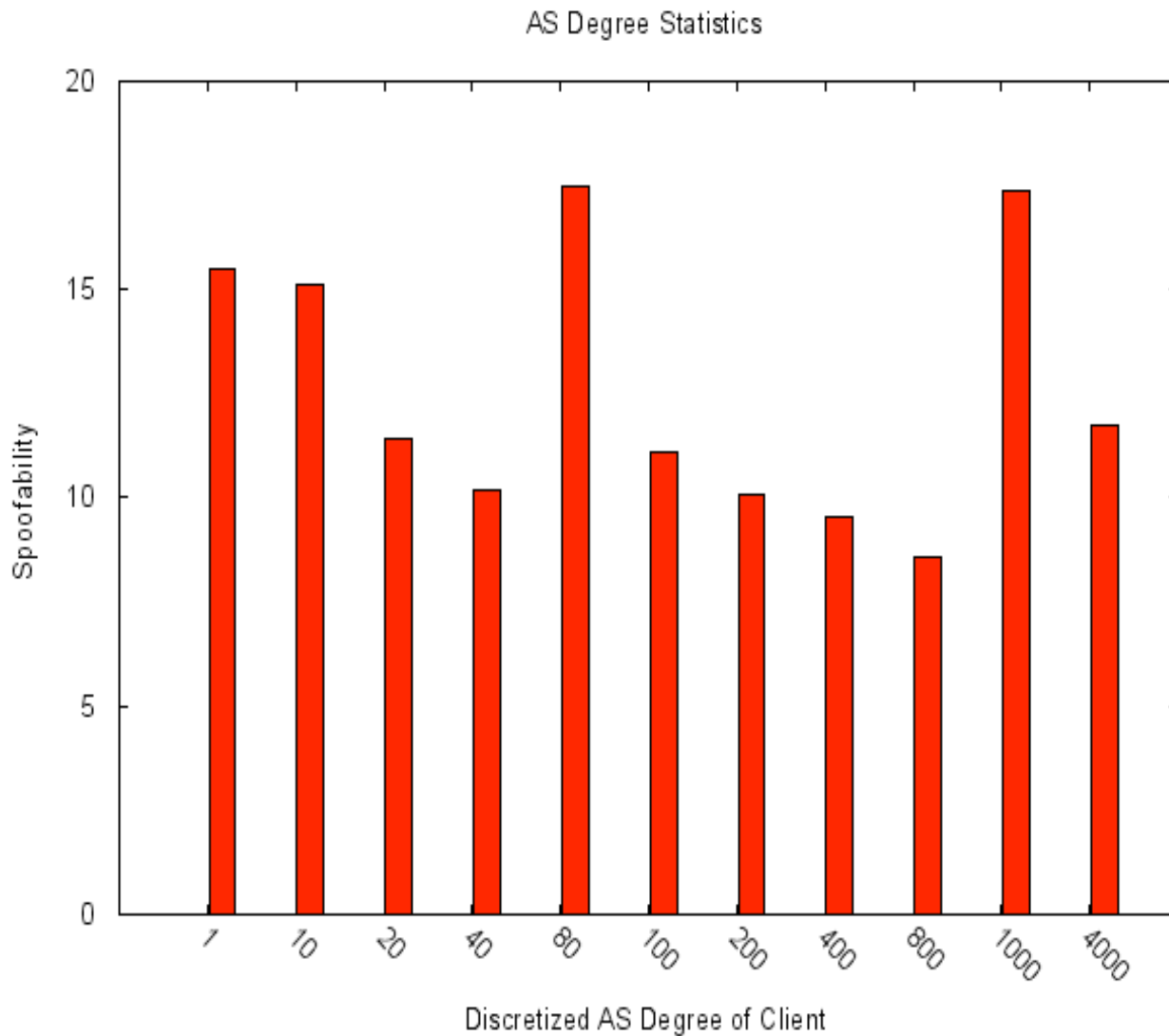
Connection Classes

Client Class Statistics



- DSL, cable easiest to control (built into architecture)
- Commercial, unknown highest spoofing rates

AS Degree



- Small or large providers filtering?
- Surprisingly, no clear trend
- Work required across the board (or a new solution)

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- What's New: Data
- Longitudinal Analysis
- Parting Thoughts

Parting Thoughts (1)

- Among clients able to spoof, what sources can they spoof?

<u>Spoofer Source</u>	<u>Description</u>	<u>Defense</u>	<u>Percent</u>
6.1.2.3	Valid (In BGP table)	uRPF	90%
1.2.3.4	Unallocated	Bogon Filters	58%
172.16.1.100	RFC1918 private	Static ACL	1%

Parting Thoughts (1)

<u>Spoofer Source</u>	<u>Description</u>	<u>Defense</u>	<u>Percent</u>
6.1.2.3	Valid (In BGP table)	uRPF	90%
1.2.3.4	Unallocated	Bogon Filters	58%
172.16.1.100	RFC1918 private	Static ACL	1%

Low hanging fruit already employed,
problem is harder!

Parting Thoughts (2)

- Tracefilter exposes operational tension between current filtering incentives and difficulty managing edge filters
- If a spoofed packet isn't filtered at edge, will travel unimpeded to destination
- Should we think about core filtering techniques?
 - StackPI
 - ML approaches with soft response (rbeverly thesis work)
 - Others

Parting Thoughts

- Even after all these years, source spoofing problem not solved
 - BCP38 has been around for 9 years
 - BCP38 great, but incentives wrong
- Single unfiltered ingress can compromise entire Internet system
 - Can we plug every hole?
 - Regulatory Response? ... but multinational?
 - Spoofer page for public provider flogging?
- What's needed (biased opinion!):
 - Clean slate design
 - Filtering in the core

Parting Thoughts

- Even after all these years, source spoofing problem not solved
 - BCP38 has been around for x years
 - BCP38 great, but incentives wrong
- Single unfiltered ingress can compromise entire Internet
 - C
 - F
 - S
- What's needed (biased opinion!):
 - Clean slate design
 - Filtering in the core

Thanks!

`http://spoofer.csail.mit.edu`